

Annihilator of a Power of a Polynomial

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22.11.2007, RWTH

Problem formulation

Given a ring $R = \mathbb{C}[x_1, \dots, x_n]$, a polynomial $f \in R$ and a number $\alpha \in \mathbb{C}$. Compute the left ideal $\text{Ann}(f^\alpha) \in D(R)$, where $D(R)$ is the Weyl algebra in $2n$ variables $x_1, \dots, x_n, \partial_1, \dots, \partial_n$ subject to usual relations.

Preliminaries

We utilize a D -module structure of a left module in

$$R[f^s] := \mathbb{C}[x_1, \dots, x_n, \frac{1}{f}] \cdot f^s.$$

The algorithm ANNFS computes a D -module structure on $R[f^s]$, that is a left ideal $I \subset D$, such that $R[f^s] \cong D/I$.

Algebraic Analysis

Indeed, for A a G -algebra of Lie type, GrA is commutative and we have $\text{GK. dim}_A(M) = \text{GK. dim}_{GrA} L(M) = \text{Kr. dim}_{GrA} L(M)$.

Theorem (Weak FTAA, SST)

A proper left ideal in $D(R) = \mathbb{K}\langle x_1, \dots, x_n, \partial_1, \dots, \partial_n \mid \dots \rangle$ has GK-dimension $\geq n$.

Let $I \subset D(R)$. Compute the left Gröbner basis of I with respect to the elimination ordering for ∂_i (a weight vector $(0, \dots, 0, 1, \dots, 1)$). Then the **characteristic ideal** of I is the ideal in the commutative ring $GrD(R) = \mathbb{K}[x_1, \dots, x_n, \partial_1, \dots, \partial_n]$, generated by the leading terms of I . The zero set of this ideal is called the **characteristic variety**.

Theorem (Strong FTAA, SST)

Let I be a proper left ideal in $D(R) = \mathbb{K}\langle x_1, \dots, x_n, \partial_1, \dots, \partial_n \mid \dots \rangle$. Then every minimal prime of a char. ideal of I has dimension $\geq n$.

Multiplicatively Closed Ore Subsets

Let S be a multiplicatively closed (**m.c.**) subset of some algebra A , that is $1 \in S$ and $a, b \in S \Rightarrow ab \in S$. S is called an **Ore set** in A , if

$$\forall s \in S, a \in A, \exists r \in S, b \in A \text{ such that } ar = sb \text{ (} s^{-1}a = br^{-1}\text{)}.$$

Ore condition

$$\forall s \in S, a \in A, \quad sA \cap aS \neq \emptyset.$$

For an associative \mathbb{K} -algebra A and a m.c. subset S , we consider $S \times A$ and introduce the following equivalence relation \simeq on it:
 $(s, a) \simeq (r, b)$, if for some $x, y \in A$, $ax = by$, $sx = ry$.

Ore localization

Then, $S \times A / \simeq$ is called an **Ore localization** of A w.r.t. S .
It is often denoted by $A_S = \{s^{-1}a \mid s \in S, a \in A\}$.

Ann F^s Method: From Kashiwara to Malgrange

Recall, that for $s \in \mathbb{K}$, $\text{Ann}_{D(R)} f^s = \{a \in D(R) \mid a \bullet f^s = 0\}$.

Theorem (Kashiwara 1981)

$D(R)/\text{Ann}_{D(R)} f^s$ is a (regular) holonomic $D(R)$ -module for any $s \in \mathbb{K}$.

Malgrange's construction for $f = f_1 \cdots f_p$: consider the left ideal

$$I_f := \left\langle \left\{ t_j - f_j, \sum_{j=1}^p \frac{\partial f_j}{\partial x_i} \partial t_j + \partial_i \right\}, 1 \leq j \leq p, 1 \leq i \leq n, \right\rangle$$

$$I_f \subset \mathbb{K}\langle \{t_j, \partial t_j\} \mid [\partial t_j, t_j] = 1 \rangle \otimes_{\mathbb{K}} \mathbb{K}\langle \{x_i, \partial_i\} \mid [\partial_i, x_i] = 1 \rangle$$

Theorem (1.)

The ideal of operators in $D[s] := D(R) \otimes_{\mathbb{K}} \mathbb{K}[s]$, annihilating f^s equals to the image of the $I_f \cap D[t \cdot \partial t]$ under the substitution $t \cdot \partial t \mapsto -s - 1$.

Proof: next slides.

Recall: Generalized Product Criterion

Let A be an associative \mathbb{K} -algebra. We use the following notations:
 $[a, b] := ab - ba$, a *commutator* or a *Lie bracket* of $a, b \in A$.

$\forall a, b, c \in A$ the following bracket identities hold

- $[a, b] = -[b, a]$, in particular $[a, a] = 0$
- $[ab, c] = a[b, c] + [a, c]b$

Recall Levandovskyy and Schönemann, ISSAC 2003.

Generalized Product Criterion

Let A be a G -algebra of Lie type (that is, all relations are of the type $x_j x_i = x_i x_j + d_{ij}$, $\forall 1 \leq i < j \leq n$).

Let $f, g \in A$. Suppose that $\text{Im}(f)$ and $\text{Im}(g)$ have no common factors, then $\text{spoly}(f, g) \rightarrow_{\{f, g\}} [f, g]$.

Proving the Theorem, part I

Let $f = f_1 \cdot \dots \cdot f_p$ and let $g_i = \partial_i + \sum_j \frac{\partial f_j}{\partial x_i} \partial t_j$. A_{p+n} below stays for $D(R) \otimes_{\mathbb{K}} \mathbb{K}\langle \{t_i, \partial t_i \mid 1 \leq i \leq p\} \mid t_i \partial t_i = \partial t_i \cdot t_i + 1 \rangle$.

Lemma

$I_f = \langle \{t_j - f_j, \{g_i\}\}, 1 \leq j \leq p, 1 \leq i \leq n \rangle \subset A_{p+n}$ is a maximal ideal, hence A_{p+n}/I_f is a simple module.

Proof.

Choose an ordering with $\{t_i, \partial_i\} \gg \{x_i, \partial t_j\}$. Running Buchberger's algorithm, we see that

$\text{spoly}(g_i, g_k) \rightarrow [g_i, g_k] = \partial t_j \sum_j [\partial_i, \frac{\partial f_j}{\partial x_k}] + \partial t_j \sum_j [\frac{\partial f_j}{\partial x_i}, \partial_k]$. Since

$[\partial_i, \frac{\partial f_j}{\partial x_k}] = \frac{\partial^2 f_j}{\partial x_i \partial x_k}$, the $\text{spoly}(g_i, g_k)$ reduces to zero.

$\text{spoly}(t_k - f_k, g_i) \rightarrow [t_k - f_k, g_i] = \sum_j \frac{\partial f_j}{\partial x_i} [t_k, \partial t_j] - [f_k, \partial_i] = 0$.

Hence, I_f is given in a Gröbner basis and its leading monomials are $\{t_j, \partial_i\}$. Thus, the GK. dim $(A/I_f) = 2(p+n) - (p+n) = p+n$, hence I_f is holonomic. □

Proving the Theorem, part II

Consider the shift algebra $\mathbb{K}\langle s, E_s \mid E_s s = s E_s + E_s = (s + 1) E_s \rangle$.

The **Mellin transform** is an injective \mathbb{K} -algebra homomorphism

$$\mathbb{K}\langle s, E_s = s E_s + E_s \rangle \rightarrow \mathbb{K}\langle t, \partial t \mid t \cdot \partial t = \partial t \cdot t + 1 \rangle, \quad s \mapsto -t \partial t - 1, \quad E_s \mapsto t.$$

Its image is the subalgebra $\mathbb{K}\langle t \partial t, t \mid \dots \rangle$.

Lemma

I_f is the annihilator of f^s in $D(R) \otimes_{\mathbb{K}} \mathbb{K}\langle \{t_j, \partial t_j\} \mid t_j \partial t_j = \partial t_j t_j + 1 \rangle$.

Proof.

The Mellin transform allows to supply $\mathbb{K}[s, \mathbf{x}, f^s]$ with the following structure of $D(R) \otimes_{\mathbb{K}} \mathbb{K}\langle t, \partial t \rangle$ -module ($p = 1$ for simplicity):

$$x_i \bullet g(s, \mathbf{x}) f^s = x_i g(s, \mathbf{x}) f^s, \quad \partial_i \bullet g(s, \mathbf{x}) f^s = \frac{\partial g}{\partial x_i} f^s + s g(s, \mathbf{x}) \frac{\partial f}{\partial x_i} f^{s-1},$$

$$t \bullet g(s, \mathbf{x}) f^s = g(s + 1, \mathbf{x}) f^{s+1}, \quad \partial t \bullet g(s, \mathbf{x}) f^s = -s g(s - 1, \mathbf{x}) f^{s-1}.$$



Proving the Theorem, part III

Proof.

In particular, for $g = 1$ and $p = 1$ we have

$$x_i \bullet f^s = x_i f^s, \quad \partial_i \bullet f^s = s \frac{\partial f}{\partial x_i} f^{s-1},$$

$$t \bullet f^s = f^{s+1}, \quad \partial t \bullet f^s = -s f^{s-1}.$$

$$\text{Then } (t_j - f_j) \bullet f^s = 0, \quad \left(\sum_{j=1}^p \frac{\partial f_j}{\partial x_i} \partial t_j + \partial_i \right) \bullet f^s = 0.$$

Since by Lemma before I_f is a maximal ideal, it is the $\text{Ann } f^s$. □

Ann F^s Algorithm in D -module Theory

Let $f = f_1 \cdot \dots \cdot f_p$.

The Ann F^s Algorithm, step I

Compute the preimage of the left ideal

$$L = \langle \{ t_j - f_j, \sum_{j=1}^p \frac{\partial f_j}{\partial x_i} \partial t_j + \partial_i \} \rangle, 1 \leq j \leq p, 1 \leq i \leq n$$

in the subalgebra $\mathbb{K} \langle \{ t_j \cdot \partial t_j \} \rangle \langle \{ x_i, \partial_i \mid [\partial_i, x_i] = 1 \} \rangle$ of

$$\mathbb{K} \langle \{ t_j, \partial t_j \} \mid [\partial t_j, t_j] = 1 \rangle \otimes_{\mathbb{K}} \mathbb{K} \langle \{ x_i, \partial_i \} \mid [\partial_i, x_i] = 1 \rangle$$

Moreover, in the preimage, $t_j \cdot \partial t_j$ will be replaced by $-s_j - 1$ (algebraic Mellin transform), where s_j are new variables, commuting with $\{x_k, \partial_k\}$.

Ann F^s Algorithm in D -module Theory

Remember, $f = f_1 \cdot \dots \cdot f_p$.

The Ann F^s Algorithm, step II

Denote the result of step I by $L' \in \mathbb{K}[\{s_j\}]\langle\{x_i, \partial x_i \mid [\partial x_i, x_i] = 1\}\rangle$. Compute the preimage of the left ideal $\langle L', f \rangle$ in the commutative subalgebra $\mathbb{K}[\{s_j\}]$.

If $p = 1$, e.g. $f = f_1$, the output is a principal ideal. Its monic generator is called a **global Bernstein–Sato polynomial** $b(s)$.

There exists an operator $B \in D(R)$, such that $B \bullet f^{s+1} = b(s) \cdot f^s$.

Theorem (Kashiwara)

All roots of $b(s)$ are rational numbers.

Note, that $s + 1$ always divides $b(s)$.

OT method for Step I

Oaku–Takayama method, 1999

$\{u_j, v_j, s_j\}$ commute with everything, $\{[\partial_i, x_i] = 1, [\partial t_j, t_j] = 1\}$.

$$\mathbb{K}\langle t_j, \partial t_j, x_i, \partial_i, u_j, v_j \mid \dots \rangle \supset \langle \{t_j - u_j f_j, \sum_{k=1}^p \frac{\partial f_k}{\partial x_i} u_k \partial t_j + \partial_i, u_j v_j - 1\} \rangle$$

1. Intersect the ideal with the subalgebra $\mathbb{K}\langle t_j, \partial t_j, x_i, \partial_i \mid \dots \rangle$ i.e. eliminate $\{u_j, v_j\}$.

2. Intersect the result of p.1. with $\mathbb{K}[-t_j \partial t_j - 1] \otimes_{\mathbb{K}} \mathbb{K}\langle x_i, \partial_i \mid \dots \rangle$, replace $-t_j \partial t_j - 1$ by s_j .

$p = 1$

$$\langle t - uf, uv - 1, \left\{ \frac{\partial f}{\partial x_i} u \partial t + \partial_i \right\} \rangle$$

The total result lives in $\mathbb{K}\langle x_i, \partial_i \mid \dots \rangle \otimes_{\mathbb{K}} \mathbb{K}[\{s\}]$

Given $\alpha \in \mathbb{C}$. Then we have the following:

Theorem (2.)

Let $f \in R = \mathbb{C}[x_1, \dots, x_n]$ and α_0 is the minimal integer root of the global b -function $b(s)$ of f . If $\alpha \notin \alpha_0 + 1 + \mathbb{N}_0$, then

$$\text{Ann}_{D(R)} f^\alpha = \text{Ann}_{D(R)[s]} f^s \Big|_{s=\alpha}.$$

For the case, when $\alpha \in \alpha_0 + 1 + \mathbb{N}_0$, we apply the **Algorithm 3**:

- 1 Compute $\text{Ann } f^s = \{g_1(s), \dots, g_r(s)\} \subset D(R)[s]$.
- 2 Compute $b(s)$ of f ; let $\alpha_0 :=$ the minimal integer root of $b(s)$.
- 3 Let $d = \alpha - \alpha_0$. If $d \leq 0$ output $\{g_i(\alpha)\}$ and stop.
- 4 For $d \in \mathbb{Z}_+$, compute the generators $\{s^{(k)}\}$ of the module

$$\text{syz}(\{f^d, g_1(\alpha_0), \dots, g_r(\alpha_0)\}) \subset D(R)^{r+1}$$

- 5 Output $\{g_i(\alpha)\} \cup \{s_1^{(k)}\}$, where $s_1^{(k)}$ is the 1st component of $s^{(k)}$.

Examples with SINGULAR:PLURAL

① (don't forget to) start SINGULAR

② load the `dmod.lib` library by typing

```
> LIB "dmod.lib";
```

③ define a commutative ring R and a polynomial F , e.g.

```
> ring R = 0, (x,y,z), dp;  poly F = x3+y3+z3;
```

④ run the `annfsOT` routine. It returns a ring, call it, say, S

```
> def S = annfsOT(F);  setring S;
```

⑤ in the ring S (= Weyl algebra of R), there are the following computed objects:

a) an ideal LD (the desired D -module structure)

b) a list BS containing the roots (with mult's) of Bernstein poly.

⑥ If you wish to compute an s -parametric annihilator, run

```
> setring R;  def P = SannfsOT(F);  setring P;
```

⑦ in the output ring the ideal LD is the parametric D -module structure

Examples

In the same way as on the previous slide you can see how the algorithms of Brianson–Maisonobe and LOT work.

- OT: `annfsOT`, `SannfsOT`
- Brianson–Maisonobe, BM: `annfsBM`, `SannfsBM`
- LOT: `annfsLOT`, `SannfsLOT`
- Multivariate BM: `annfsBMI`
- All the relevant data at once: `operatorBM`

If you wish to see progress of each step of the algorithm, set before computation `printlevel = 1;`. If you wish to see additionally all intermediate data, set `printlevel = 2;`.

Example Session

`x`; `x4`; `x3 - y2`; `y5 + xy4 + x4` (a Reiffen curve), `x3 + y3 + z2w` (4 variables), `(x3 + y2) · (x2 + y3)`.

Non-genericity with SINGULAR:PLURAL

Assume we have $F = x^3 + y^3 + z^3$ and we'd like to compute $\text{Ann } F^n$ for $n \geq 1$. As we know from the example before, the minimal integer root is -2 . So, any $n \geq -1$ leads to the exceptional situation (Algorithm 3 instead of Theorem 2 above). Let us compute the structure of the annihilator for $n = 3$.

- 1 define a commutative ring R and a polynomial F , e.g.
> ring R = 0, (x, y, z), dp; poly F = x3+y3+z3;
- 2 run the `SannfsBM` routine. It returns a ring, call it, say, S
> def S = SannfsBM(F); setring S;
- 3 in the ring $S = D(R)[s]$, there is an ideal LD (s -parametric D -module structure)
> int n = 3; poly F = imap(R, F);
> ideal I = annfspecial(LD, F, -2, n); I = groebner(I);
- 4 the ideal I is the desired D -module structure of F^3 .

Thank you for your attention!

Please visit the SINGULAR homepage

- <http://www.singular.uni-kl.de/>
- there you find among others the online manual (with detailed documentation and examples for each command, procedure and library)

Integrals

For generic f , an integral $\int_C f(x, t)^\alpha t^\gamma dt$ satisfies a Gel'fand–Kapranov–Zelevinsky system [SST]. However, one needs to treat non-generic polynomials too.

Hypergeometric Integral

$$\text{Consider } F(\alpha; x) = \int_C \prod_{i=1}^p f_i(x, t)^{\alpha_i} dt_1 \cdots dt_m,$$

where $\alpha_i \in \mathbb{K} \subseteq \mathbb{C}$ and C is an m -cycle. The function $F(\alpha; x)$ depends on the homology class of C . Let

$$D = \mathbb{K}\langle t_1, \dots, t_m, \partial_{t_1}, \dots, \partial_{t_m}, x_1, \dots, x_n, \partial_{x_1}, \dots, \partial_{x_n} \mid \\ \partial_{t_j} t_i = t_i \partial_{t_j} + \delta_{ij}, \partial_{x_j} x_i = x_i \partial_{x_j} + \delta_{ij} \rangle$$

Theorem (SST, Th. 5.5.1)

Let $I \subset D$ be a left ideal, annihilating the function

$f_\alpha(x, t) = \prod_{i=1}^p f_i(x, t)^{\alpha_i}$, $f_i \in \mathbb{K}[x_1, \dots, x_n, t_1, \dots, t_m]$. Then, the ideal

$$J = (I + \langle \partial_{t_1}, \dots, \partial_{t_m} \rangle_D) \cap \mathbb{K}\langle x_1, \dots, x_n, \partial_{x_1}, \dots, \partial_{x_n} \mid \partial_{x_j} x_i = x_i \partial_{x_j} + \delta_{ij} \rangle$$

annihilates the function $F(\alpha; x)$.

The left ideal J is called the integral ideal of I with respect to t .

Note

In the Theorem, we have to intersect the sum of a left and a right ideals with a subalgebra. There are no general methods (only very specific, e.g. of Takayama) for treating this situation.