## Localisation of D-modules

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#### Introduction and Motivation

(1) The Annihilator (2) The Generator Conclusion

### Notation

- K computable field of characteristic 0 contained in  $\mathbb C$
- $R_n := K[\underline{x}] := K[x_1, .., x_n]$  the polynomial ring in n indeterminates
- $D_n:=R\langle\underline{\partial}
  angle:=R_n\langle\partial_1,..,\partial_n
  angle$  be the n-th Weyl algebra

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# Aim and Motivation

- Let  $f \in R_n$  and M a holonomic (left)  $D_n$ -module  $M \cong D_n/I$  for a left ideal I in  $D_n$
- Compute  $M[f^{-1}] := R_n[f^{-1}] \otimes_{R_n} M$
- Can be generalised to M only being holonomic on  $K^n \setminus \mathcal{V}(f)$
- "localise away" non-holonomic locus
- Generalisation to last weeks with  $M = R_n$

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### Plan

- ullet Want to find a generator  $f^a\otimes 1$  and its annihilator
- Call this generator  $f^s \otimes 1 \otimes 1 \in f^s \otimes_K R_n[f^{-1}, s] \otimes_{R_n} M$
- (1) Compute  $J^{I}(f^{s}) := \operatorname{Ann}_{D_{n}[s]}(f^{s} \otimes 1 \otimes 1)$
- (2) Compute a suitable number  $a \in K$  for substituting s by a
  - $\underline{x}$  operates by left multiplication on the right factor

• 
$$\partial_i \bullet (f^s \otimes \frac{g(x,s)}{f^k} \otimes Q) =$$
  
 $f^s \otimes \frac{sg(x,s)f_i}{f^{k+1}} \otimes Q + f^s \otimes \partial_i(\frac{g(x,s)}{f^k}) \otimes Q + f^s \otimes \frac{g(x,s)}{f^k} \otimes \partial_i Q$   
•  $f_i := \frac{\partial f}{\partial x_i}.$ 

Idea Getting many Generators Intersecting away

# (1) The Annihilator

Aim of this section: Compute  $J^I(f^s):= {\sf Ann}_{D_n[s]}(f^s\otimes 1\otimes 1)$ 

Idea Getting many Generators Intersecting away

# Idea

• Extend 
$$D_n[s]$$
 to  $D_{n+1} := D_n \langle t, \partial_t \rangle$   
•  $t \bullet (f^s \otimes \frac{g(x,s)}{f^k} \otimes Q) := f^s \otimes \frac{g(x,s+1)f}{f^k} \otimes Q$   
•  $\partial_t \bullet (f^s \otimes \frac{g(x,s)}{f^k} \otimes Q) := f^s \otimes \frac{-sg(x,s-1)}{f^{k+1}} \otimes Q$   
• Try to compute  $J_{n+1}^I(f^s) := \operatorname{Ann}_{D_{n+1}}(f^s \otimes 1 \otimes 1)$   
• "Intersect" this with  $D_n[s]$ 

• 
$$-\partial_t t$$
 acts by  $s$ , so  $D_n[s] \hookrightarrow D_{n+1}$ 

Idea Getting many Generators Intersecting away

# Getting many Generators

• 
$$\phi: D_{n+1} \xrightarrow{\sim} D_{n+1}: x_i \mapsto x_i, t \mapsto t - f, \partial_i \mapsto \partial_i + f_i \partial_t, \partial_t \mapsto \partial_t$$
  
• Lemma: Let  $I$  be  $f$ -saturated. Then

$$J_{n+1}^{I}(f^{s}) =_{D_{n+1}} \langle \phi(I), t - f \rangle$$

holds.

Idea Getting many Generators Intersecting away

## Intersecting away

- Input: Left ideal I of  $D_{n+1}$
- Output:  $J = I \cap D_n[s] = I \cap D_n[-\partial_t t]$
- Weight vector w on  $D_{n+1}[y_1, y_2]$  by  $w(t) = 1, w(\partial_t) = -1, w(x_i) = w(\partial_i) = 0, w(y_1) = 1, w(y_2) = -1$
- Homogenize I by  $y_1$  according to w
- Compute Gröbner basis  $\tilde{J}$  of this ideal and  $1-y_1y_2$  eliminating  $y_1$  and  $y_2$
- Take elements of  $\tilde{J}$  not having  $y_1$  or  $y_2$  and multiply them with appropriate powers of t and  $\partial_t$  to give them a w-degree of 0
- Return these elements
- The above lemma in combination with this algorithm gives the solution to this section's problem.

Bernstein Polynomial Determine the Exponent



Aim of this section:

Compute a suitable number  $a \in K$  to get a generator and to substitute s by a in last section's result

Bernstein Polynomial Determine the Exponent

# Bernstein Polynomial

• Bernstein polynomial  $b_f^I(s) \in K[s]$ : the monic generator for all  $b \in K[s]$ , s.t. exists  $Q(s) \in D_n[s]$  with:

$$b(s) \bullet (f^s \otimes 1 \otimes 1) = Q(s) \bullet (f^s \otimes f \otimes 1) = Q(s) f \bullet (f^s \otimes 1 \otimes 1)$$

Fix Q<sup>I</sup><sub>f</sub>(s) as operator with above properties
Idea: Q<sup>I</sup><sub>f</sub>(s) is some kind of "inverse" for f

Bernstein Polynomial Determine the Exponent

# Computing the Bernstein Polynomial

- Input:  $f \in R_n$  and f-saturated holonomic ideal  $I \trianglelefteq D_n$
- Output:  $b_f^I(s)$
- Compute  $J^{I}(f^{s})$  by means of section 1
- Compute the monic generator of  $_{D_n[s]}\langle f, J^I(f^s)\rangle \cap K[s]$

Bernstein Polynomial Determine the Exponent

#### Determine the Exponent

Theorem: M = D<sub>n</sub>/I holonomic and a ∈ K<sup>\*</sup>, such that no element of {a − 1, a − 2, ..} is root of b<sup>I</sup><sub>f</sub>(s), then:

 $f^{a} \otimes_{K} R_{n}[f^{-1}] \otimes_{R_{n}} M$  $\cong D_{n} \bullet (f^{a} \otimes 1 \otimes 1)$  $\cong (D_{n}[s]/J^{I}(f^{s}))|_{s=a}$ 

• Take a as smallest negative integer root of  $b_f^I(s)$ . If no such number exists, then a := -1

# Final Algorithm

- $\bullet$  Input:  $f\in R_n,\ M=D_n/I$  holonomic and  $f\mbox{-saturated}$
- Output:  $J \leq D_n$  and  $a \in \mathbb{Z}$  with  $R_n[f^{-1}] \otimes_{R_n} M \cong D_n/J$  generated by  $f^a \otimes 1$ .
- Determine  $J^{I}(f^{s})$  as in section 1
- Determine  $b_f^I(s)$  as in section 2
- Find the smallest integer root a of  $b_f^I(s)$ . If not exist, a:=-1
- Replace s by a in each generator of  $J^{I}(f^{s}) \rightsquigarrow J$

### Final words:

