

Lattices and Codes, analogies and interactions. or Codes, Invariants and Modular Forms, a conclusion. Gabriele Nebe, RWTH Aachen Bonn, 11.7.2008 Codes.

*R* finite ring, *A* finite left *R*-module  $C \le A^N$  is called a **code**, *N* its **length** and **codewords**  $c = (c_1, \ldots, c_N)$  are rows. The **codepolynomial** of *C* is

$$p_C := \sum_{c \in C} \prod_{i=1}^N x_{c_i} \in \mathbb{C}[x_a \mid a \in A]_N = p_C^{(1)}$$

The genus m codepolynomial of C is

$$p_C^{(m)} := \sum_{(c^{(1)}, \dots, c^{(m)}) \in C^m} \prod_{v \in A^m} x_v^{a_v(c^{(1)}, \dots, c^{(m)})} \in \mathbb{C}[x_v : v \in A^m].$$

where

$$a_v(c^{(1)}, \dots, c^{(m)}) := |\{j \in \{1, \dots, N\} \mid c_j^{(i)} = v_i \text{ for all } i \in \{1, \dots, m\}\}|$$
  
for  $v := (v_1, \dots, v_m) \in A^m$ .

For  $C \leq A^N$  and  $m \in \mathbb{N}$  let  $C(m) := R^{m \times 1} \otimes C = \{(c^{(1)}, \dots, c^{(m)})^{\mathsf{Tr}} \mid c^{(1)}, \dots, c^{(m)} \in C\} \leq (A^m)^N$ Then

$$p_C^{(m)} = p_{C(m)}.$$

A typical element of C(m) is a matrix in  $A^{m \times N}$ , where the rows are codewords in C.



**The finite Siegel**  $\Phi$ **-operator.** (B. Runge, 1995)

$$\Phi_m: p_C^{(m)} \mapsto p_C^{(m-1)}$$

is given by the variable substitution:

$$x_{(v_1,\dots,v_m)} \mapsto \begin{cases} x_{(v_1,\dots,v_{m-1})} & \text{if } v_m = 0\\ 0 & \text{else} \end{cases}$$

 $p_C^{(m-1)}$  is obtained from  $p_C^{(m)}$  by counting only those matrices

in which the last row is zero.

Lattices and Theta Series.

 $L \leq (\mathbb{R}^N, (,))$  a lattice in Euclidean N-space.

The **theta series** of L is

$$\vartheta_L(z) = \sum_{\ell \in L} q^{(\ell,\ell)}$$

where  $q = \exp(\pi i z)$ .

The genus m Siegel theta series of L is

$$\vartheta_L^{(m)}(Z) = \sum_{\underline{\ell} \in L^m} \exp(\pi i \operatorname{Tr}(Z(\underline{\ell}, \underline{\ell}))).$$

The Siegel  $\Phi$ -operator maps  $\vartheta_L^{(m)}$  to  $\vartheta_L^{(m-1)}$ .

#### Codes and Lattices: Construction A.

Let p be a prime and  $(b_1,\ldots,b_N)$  be a basis of  $\mathbb{R}^N$  such that

$$(b_i, b_j) = \begin{cases} 0 & \text{if } i \neq j \\ 1/p & \text{if } i = j \end{cases}$$

Let  $C \leq \mathbb{F}_p^N = \mathbb{Z}^N / p\mathbb{Z}^N$  be a code. Then the **codelattice**  $L_C$  is  $L_C := \{\sum_{i=1}^N a_i b_i \mid (a_1 \pmod{p}), \dots, a_N \pmod{p}) \in C\}$ 

#### Remark.

(a)  $L_C^* = L_{C^{\perp}}$ , so  $L_C$  is unimodular, if C is self-dual. (b)  $L_C$  is even unimodular, if p = 2 and C is a Type II code. (c)  $\vartheta_{L_C} = p_C(\vartheta_0, \dots, \vartheta_{p-1})$  where

$$\vartheta_a = \vartheta_{(a+p\mathbb{Z})b_1} = \sum_{n=-\infty}^{\infty} q^{(a+pn)^2/p}$$

similarly for higher genus theta series and codepolynomials.

#### Theta series are Modular Forms.

If  $L = L^*$  and  $(\ell, \ell) \in 2\mathbb{Z}$  for all  $\ell \in L$ , even unimodular lattice, then

$$\vartheta_L^{(m)}(Z) = \sum_{\underline{\ell} \in L^m} \exp(\pi i \operatorname{Tr}(Z(\underline{\ell}, \underline{\ell}))) \in \mathcal{M}_{N/2}(\operatorname{Sp}_{2m}(\mathbb{Z}))$$

where

$$\operatorname{Sp}_{2m}(\mathbb{Z}) = \langle \begin{pmatrix} A & 0 \\ 0 & A^{-tr} \end{pmatrix}, \begin{pmatrix} I_m & B \\ 0 & I_m \end{pmatrix}, \begin{pmatrix} 0 & -I_m \\ I_m & 0 \end{pmatrix} \\ | A \in \operatorname{GL}_m(\mathbb{Z}), B = B^{tr} \in \mathbb{Z}^{m \times m} \rangle.$$

### Codepolynomials are Invariants.

R finite ring, A finite left R-module,  $\beta:A\times A\to \mathbb{Q}/\mathbb{Z}$  regular. For  $C\leq A^N$  the dual code is

$$C^{\perp} := \{ v \in A^N \mid \sum_{i=1}^N \beta(v_i, c_i) = 0 \text{ for all } c \in C \}.$$

Let  $M := \{\beta^r : (x, y) \mapsto \beta(x, ry) \mid r \in R\}$  and assume that  $M \cong R_R$ and is closed under interchanging arguments.

Additional quadratic conditions are given by a subgroup  $Q \leq (\mathbb{Q}/\mathbb{Z})^A$ , such that:

- For all  $\varphi \in Q$ ,  $\lambda(\varphi) : (x, y) \mapsto \varphi(x + y) \varphi(x) \varphi(y) \in M$ .
- For all  $r \in R$ ,  $\varphi \in Q$ ,  $\varphi[r] : x \mapsto \varphi(rx) \in Q$ .
- For all  $r \in R$ ,  $\{\!\!\{ \beta^r \}\!\!\} : x \mapsto \beta(x, rx) \in Q$ . Then  $(R, A, \beta, Q)$  is called a **Type**.

C is called a **Type** T **code**, if  
a) 
$$C \leq A^N$$
 is an *R*-module.  
b)  $\sum_{i=1}^{N} \varphi(c_i) = 0$  for all  $\varphi \in Q$ ,  $c \in C$  (isotropic).  
c)  $C = C^{\perp}$  (self-dual)

Examples. Type I codes (2<sub>I</sub>)

$$R = \mathbb{F}_2 = A, \ \beta(x, y) = \frac{1}{2}xy, \ Q = \{\varphi : x \mapsto \frac{1}{2}x^2 = \beta(x, x), 0\}$$

Type II code  $(2_{II})$ .

$$R = \mathbb{F}_2 = A, \ \beta(x, y) = \frac{1}{2}xy, \ Q = \{\phi : x \mapsto \frac{1}{4}x^2, 2\phi = \varphi, 3\phi, 0\}$$

Type IV codes  $(4^H)$ .

$$R = \mathbb{F}_4 = A, \ \beta(x, y) = \frac{1}{2} \operatorname{trace}(x\overline{y}), \ Q = \{\varphi : x \mapsto \frac{1}{2}x\overline{x}, 0\}$$
  
where  $\overline{x} = x^2$ .

Additive codes over  $\mathbb{F}_4$ . (4<sup>*H*+</sup>)

$$R = \mathbb{F}_2, \ A = \mathbb{F}_4, \ \beta(x, y) = \frac{1}{2} \operatorname{trace}(x\overline{y}), \ Q = \{\varphi : x \mapsto \frac{1}{2}x\overline{x}, 0\}$$

#### Clifford-Weil groups.

Let  $T := (R, A, \beta, Q)$  be a Type. Then the associated Clifford-Weil group C(T) is a subgroup of  $GL_{|A|}(\mathbb{C})$ 

 $C(T) = \langle m_r, d_{\varphi}, h_{e,u_e,v_e} | r \in R^*, \varphi \in Q, e = u_e v_e \in R$  symmetric idempotent  $\rangle$ Let  $(x_a | a \in A)$  denote a basis of  $\mathbb{C}^{|A|}$ . Then

 $m_r : x_a \mapsto x_{ra}, \quad d_{\varphi} : x_a \mapsto \exp(2\pi i \varphi(a)) x_a$ 

$$h_{e,u_e,v_e}$$
:  $x_a \mapsto |eA|^{-1/2} \sum_{b \in eA} \exp(2\pi i\beta(b,v_ea)) x_{b+(1-e)a}$ 

Similarly the genus m Clifford-Weil group

 $\mathcal{C}_m(T) = \langle m_r, d_{\varphi}, h_{e,u_e,v_e} \mid r \in \mathsf{GL}_m(R), \varphi \in Q^{(m)}, e = u_e v_e \in R^{m \times m} \text{ sym. id. } \rangle$  $\leq \mathsf{GL}_{|A|^m}(\mathbb{C})$ 

$$m_r : x_a \mapsto x_{ra}, \quad d_{\varphi} : x_a \mapsto \exp(2\pi i \varphi(a)) x_a$$
  
 $h_{e,u_e,v_e} : x_a \mapsto |eA|^{-1/2} \sum_{b \in eA} \exp(2\pi i \beta(b, v_e a)) x_{b+(1-e)a}$ 

## Theorem.

Let  $C \leq A^N$  be a self-dual isotropic code of Type T. Then  $p_C^{(m)}$  is invariant under  $\mathcal{C}_m(T)$ .

## Proof.

Invariance under  $m_r$   $(r \in GL_m(R))$  because C is a code. Invariance under  $d_{\varphi}$   $(\varphi \in Q^{(m)})$  because C is isotropic. Invariance under  $h_{e,u_e,v_e}$  because C is self dual.

The main theorem. (N, Rains, Sloane (1999-2006)) If R is a direct product of matrix rings over chain rings, then

$$Inv(\mathcal{C}_m(T)) = \langle p_C^{(m)} | C \text{ of Type } T \rangle.$$

**Example:**  $C_2(II)$ .

$$R = \mathbb{F}_2^{2 \times 2}, R^* = \mathrm{GL}_2(\mathbb{F}_2) = \langle a := \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}, \ b := \begin{pmatrix} 0 & 1 \\ 1 & 1 \end{pmatrix} \rangle$$

 $A = \mathbb{F}_2^2 = \{ \begin{pmatrix} 0 \\ 0 \end{pmatrix}, \begin{pmatrix} 1 \\ 0 \end{pmatrix}, \begin{pmatrix} 0 \\ 1 \end{pmatrix}, \begin{pmatrix} 1 \\ 1 \end{pmatrix} \}, \text{symmetric idempotent } e = \text{diag}(1,0)$ 

•

$$\mathcal{C}_{2}(\mathrm{II}) = \langle m_{a} = \begin{pmatrix} 1000\\0010\\0100\\0001 \end{pmatrix}, \ m_{b} = \begin{pmatrix} 1000\\0001\\0001\\0100\\0010 \end{pmatrix}, h_{e,e,e} = \frac{1}{\sqrt{2}} \begin{pmatrix} 1100\\1-100\\0011\\001-1 \end{pmatrix}, \ d_{\phi e} = \mathrm{diag}(1, i, 1, i) \rangle$$

 $\mathcal{C}_2(II)$  has order 92160 and Molien series

$$\frac{1+t^{32}}{(1-t^8)(1-t^{24})^2(1-t^{40})}$$

where the generators correspond to the genus 2 codepolynomials of the codes:

$$e_8, g_{24}, d_{24}^+, d_{40}^+, \text{ and } d_{32}^+$$

 $\mathcal{C}_2(II)$  has a reflection subgroup of index 2, No. 31 on the Shephard-Todd list.

Higher genus Clifford-Weil groups for the classical Types of codes over finite fields.

 $\mathcal{C}_m(T) = S.(\ker(\lambda) \times \ker(\lambda)).\mathcal{G}_m(T)$ 

 $\lambda(\varphi): (x,y) \mapsto \varphi(x+y) - \varphi(x) - \varphi(y)$ 

R	J	$\epsilon$	$\mathcal{G}_m(T)$
$\mathbb{F}_q\oplus\mathbb{F}_q$	$(r,s)^J = (s,r)$	1	$GL_{2m}(\mathbb{F}_q)$
$\mathbb{F}_{q^2}$	$r^J = r^q$	1	$U_{2m}(\mathbb{F}_{q^2})$
$\mathbb{F}_q, \; q \; odd$	$r^J = r$	1	$Sp_{2m}(\mathbb{F}_q)$
$\mathbb{F}_q, \; q \; odd$	$r^J = r$	-1	$O_{2m}^+(\mathbb{F}_q)$
$\mathbb{F}_q, \; q \; even$	doubly even	•	$Sp_{2m}(\mathbb{F}_q)$
$\mathbb{F}_q, \; q \; even$	singly even		$O_{2m}^+(\mathbb{F}_q)$

## Hecke operators for codes.

## Motivation.

Determine linear relations between  $p_C^{(m)}$  for  $C \in M_N(T) = \{C = C^{\perp} \leq A^N \mid C \text{ isotropic } \}.$ 

 $M_{16}(\text{II}) = [e_8 \perp e_8] \cup [d_{16}^+]$  and these two codes have the same genus 1 and 2 codepolynomials, but  $p^{(3)}(e_8 \perp e_8)$  and  $p^{(3)}(d_{16}^+)$  are linearly independent.

 $h(M_{24}(\text{II})) = 9$  and only the genus 6 codepolynomials are linearly independent, there is one relation for the genus 5 codepolynomials.

 $h(M_{32}(\text{II})) = 85$  and here the genus 10 codepolynomials are linearly independent, whereas there is a unique relation for the genus 9 codepolynomials.

Three different approaches:

1) Determine all the codes and their codepolynomials. If dim(C) = n = N/2 there are  $\prod_{i=0}^{d-1} (2^n - 2^i)/(2^d - 2^i)$  subspaces of dimension d in C. N = 32, d = 10 yields more than  $10^{18}$  subspaces.

2) Use Molien's theorem:  $Inv_N(\mathcal{C}_m(II)) = \langle p_C^{(m)} | C \in M_N(II) \rangle$ and if  $a_N := dim(Inv_N(\mathcal{C}_m(II)))$  then

$$\sum_{N=0}^{\infty} a_N t^N = \frac{1}{|\mathcal{C}_m(\mathrm{II})|} \sum_{g \in \mathcal{C}_m(\mathrm{II})} (\det(1 - tg))^{-1}$$

Problem:  $C_{10}(II) \leq GL_{1024}(\mathbb{C})$  has order  $> 10^{69}$ .

3) Use Hecke operators.

Fix a Type  $T = (\mathbb{F}_q, \mathbb{F}_q, \beta, Q)$  of self-dual codes over a finite **field** with q elements.

$$M_N(T) = \{ C = C^{\perp} \leq \mathbb{F}_q^N \mid C \text{ isotropic } \} = [C_1] \cup \ldots \cup [C_h]$$

where [C] denotes the **permutation equivalence** class of the code C. Then  $n := \frac{N}{2} = \dim(C)$  for all  $C \in M_N(T)$ .  $C, D \in M_N(T)$  are called **neighbours**, if  $\dim(C) - \dim(C \cap D) = 1$ ,  $C \sim D$ .

$$\mathcal{V} = \mathbb{C}[C_1] \oplus \ldots \oplus \mathbb{C}[C_h] \cong \mathbb{C}^h$$

$$K_N(T) \in \mathsf{End}(\mathcal{V}), \ K_N(T) : [C] \mapsto \sum_{D \in M_N(T), D \sim C} [D].$$

#### Kneser-Hecke operator.

(adjacency matrix of neighbouring graph)

**Example.**  $M_{16}(II) = [e_8 \perp e_8] \cup [d_{16}^+]$ 



$$K_{16}(\text{II}) = \left(\begin{array}{cc} 78 & 49\\ 70 & 57 \end{array}\right)$$

 $\ensuremath{\mathcal{V}}$  has a Hermitian positive definite inner product defined by

 $\langle [C_i], [C_j] \rangle := |\operatorname{Aut}(C_i)| \delta_{ij}.$ 

**Theorem.** (N. 2006)

The Kneser-Hecke operator K is a self-adjoint linear operator.

$$\langle v, Kw \rangle = \langle Kv, w \rangle$$
 for all  $v, w \in \mathcal{V}$ .

Example. 
$$\frac{7}{10} = \frac{|\operatorname{Aut}(e_8 \perp e_8)|}{|\operatorname{Aut}(d_{16}^+)|} = \frac{49}{70}$$
 hence  
diag $(7, 10)K_{16}(\operatorname{II})^{\mathsf{Tr}} = K_{16}(\operatorname{II})$  diag $(7, 10)$ .



$$p^{(m)}: \mathcal{V} \to \mathbb{C}[X], \sum_{i=1}^{h} a_i[C_i] \mapsto \sum_{i=1}^{h} a_i p_{C_i}^{(m)}$$

is a linear mapping with kernel

$$\mathcal{V}_m := \ker(p^{(m)}).$$

Then

$$\mathcal{V} =: \mathcal{V}_{-1} \geq \mathcal{V}_0 \geq \mathcal{V}_1 \geq \ldots \geq \mathcal{V}_n = \{0\}.$$

is a filtration of  $\ensuremath{\mathcal{V}}$  yielding the orthogonal decomposition

$$\mathcal{V} = \bigoplus_{m=0}^{n} \mathcal{Y}_{m}$$
 where  $\mathcal{Y}_{m} = \mathcal{V}_{m-1} \cap \mathcal{V}_{m}^{\perp}$ .

$$\mathcal{V}_0 = \{\sum_{i=1}^h a_i [C_i] \mid \sum_{i=1}^h a_i = 0\}$$

and

$$\mathcal{V}_0^{\perp} = \mathcal{Y}_0 = \langle \sum_{i=1}^h \frac{1}{|\operatorname{Aut}(C_i)|} [C_i] \rangle.$$

# **Theorem.** (N. 2006) The space $\mathcal{Y}_m = \mathcal{Y}_m(N)$ is the $K_N(T)$ -eigenspace to the eigenvalue $\nu_N^{(m)}(T)$ with $\nu_N^{(m)}(T) > \nu_N^{(m+1)}(T)$ for all m.

Туре	$\nu_N^{(m)}(T)$
$q_{\mathbf{I}}^{E}$	$(q^{n-m}-q-q^m+1)/(q-1)$
$q_{\mathbf{II}}^E$	$(q^{n-m-1}-q^m)/(q-1)$
$q^E$	$(q^{n-m}-q^m)/(q-1)$
$q_1^E$	$(q^{n-m-1}-q^m)/(q-1)$
$q^H$	$(q^{n-m+1/2}-q^m-q^{1/2}+1)/(q-1)$
$q_1^H$	$(q^{n-m-1/2}-q^m-q^{1/2}+1)/(q-1)$

**Corollary.** The neighbouring graph is connected. Proof. The maximal eigenvalue  $\nu_0$  of the adjacency matrix is simple with eigenspace  $\mathcal{Y}_0$ .

Example: 
$$M_{16}(II) = [e_8 \perp e_8] \cup [d_{16}^+]$$
  
 $(2^{8-m-1} - 2^m : m = 0, 1, 2, 3) = (127, 62, 28, 8)$ 

$$K_{16}(\mathrm{II}) = \left(\begin{array}{cc} 78 & 49\\ 70 & 57 \end{array}\right)$$

has eigenvalues 127 and 8 with eigenvectors (7, 10) and (1, -1). Hence

$$\mathcal{Y}_0 = \langle 7[e_8 \perp e_8] + 10[d_{16}^+] \rangle$$
$$\mathcal{Y}_1 = \mathcal{Y}_2 = 0$$
$$\mathcal{Y}_3 = \langle [e_8 \perp e_8] - [d_{16}^+] \rangle.$$

Even unimodular lattices.

$$\mathcal{L}_N = \{L = L^* \leq \mathbb{R}^N \mid L \text{ even } \} = [L_1] \cup \ldots \cup [L_h]$$

where [L] denotes the **isometry** class of the lattice L.  $L, M \in \mathcal{L}_N$  are called **p-neighbours**, if  $[L : L \cap M] = p$ , notation:  $L \sim M$ .

$$\mathcal{V} = \mathbb{C}[L_1] \oplus \ldots \oplus \mathbb{C}[L_h] \cong \mathbb{C}^h$$
$$K_{N/2}(p) \in \mathsf{End}(\mathcal{V}), \ K_{N/2}(p) : [L] \mapsto \sum_{M \in \mathcal{L}_N, M \sim L} [M].$$

## Kneser-Hecke operator.

(adjacency matrix of neighbouring graph)

**Example.**  $\mathcal{L}_{16} = [E_8 \perp E_8] \cup [D_{16}^+]$ 



$$K_8(2) = \left(\begin{array}{cc} 14670 & 18225\\ 12870 & 20025 \end{array}\right)$$

 ${\cal V}$  has a Hermitian positive definite inner product defined by

 $\langle [L_i], [L_j] \rangle := |\operatorname{Aut}(L_i)| \delta_{ij}.$ 

Theorem. (Venkov, N. 2001)

The Kneser-Hecke operator K is a self-adjoint linear operator.

$$\langle v, Kw \rangle = \langle Kv, w \rangle$$
 for all  $v, w \in \mathcal{V}$ .

Example. 
$$\frac{405}{286} = \frac{|\operatorname{Aut}(E_8 \perp E_8)|}{|\operatorname{Aut}(D_{16}^+)|} = \frac{18225}{12870}$$
 hence  
diag(405, 286) $K_8(2)^{\mathsf{Tr}} = K_8(2)$  diag(405, 286).



$$\vartheta^{(m)}: \mathcal{V} \to \mathcal{M}_{N/2}(\mathsf{Sp}_{2m}(\mathbb{Z})), \sum_{i=1}^{h} a_i[L_i] \mapsto \sum_{i=1}^{h} a_i \vartheta_{L_i}^{(m)}$$

is a linear mapping with kernel

$$\mathcal{V}_m := \ker(\vartheta^{(m)}).$$

Then

$$\mathcal{V} =: \mathcal{V}_{-1} \geq \mathcal{V}_0 \geq \mathcal{V}_1 \geq \ldots \geq \mathcal{V}_N = \{0\}.$$

is a filtration of  $\ensuremath{\mathcal{V}}$  yielding the orthogonal decomposition

$$\star_L \quad \mathcal{V} = \bigoplus_{m=0}^N \mathcal{Y}_m \text{ where } \mathcal{Y}_m = \mathcal{V}_{m-1} \cap \mathcal{V}_m^{\perp}.$$

$$\mathcal{V}_0 = \{\sum_{i=1}^h a_i[L_i] \mid \sum a_i = 0\} \text{ and } \mathcal{V}_0^\perp = \mathcal{Y}_0 = \langle \sum_{i=1}^h \frac{1}{|\operatorname{Aut}(L_i)|} [L_i] \rangle.$$

**Theorem.**  $\star_L$  is invariant under  $K_{N/2}(p)$  (but the eigenspace decomposition is usually much finer and I do not know how to predict eigenvalues).

**Example:**  $\mathcal{L}_{16} = [E_8 \perp E_8] \cup [D_{16}^+]$ 

$$K_8(2) = \left(\begin{array}{cc} 14670 & 18225\\ 12870 & 20025 \end{array}\right)$$

has eigenvalues 32895 and 1800 with eigenvectors (286,405) and (1,-1). Here  $\vartheta^{(m)} = p^{(m)}(\vartheta_a : a \in \mathbb{F}_2^m)$  and all lattices come from codes.  $\mathcal{Y}_0 = \langle 286[E_8 \perp E_8] + 405[D_{16}^+] \rangle$  $\mathcal{Y}_1 = \mathcal{Y}_2 = \mathcal{Y}_3 = 0$  $\mathcal{Y}_4 = \langle [E_8 \perp E_8] - [D_{16}^+] \rangle.$  Dimension 24: The 24 Niemeier lattices. (N, Venkov)

Here h = 24 and only 9 of the lattices are codelattices. With B. Venkov we calculated  $K_{12}(2)$  and its eigenspace decomposition.

j	0	1	2	3	4	5	6	7	8	9	10	11	12
$\dim(\mathcal{Y}_j)$	1	1	1	1	2	2	3	3*	4*	2*	2*	1	1

\* means that the dimension is only conjectured.

Hecke operators as double cosets. (Lattices.)

$$t_m^{(m-1)}(p^2) := \operatorname{diag}(1, \underbrace{p, \dots, p}_{m-1}, p^2, \underbrace{p, \dots, p}_{m-1}) \in \operatorname{GSp}_{2m}(\mathbb{Z})$$

Then the double coset

$$T_m^{(m-1)}(p^2) := \operatorname{Sp}_{2m}(\mathbb{Z}) t_m^{(m-1)}(p^2) \operatorname{Sp}_{2m}(\mathbb{Z}) = \bigcup_{j=1}^d \operatorname{Sp}_{2m}(\mathbb{Z}) \gamma_j$$

acts on the space of modular forms  $\mathcal{M}_k(\operatorname{Sp}_{2m}(\mathbb{Z}))$  and also on the subspace spanned by theta series by

$$\delta_k(T_m^{(m-1)}(p^2)) : f \mapsto \sum_{j=1}^d f_{|_k \gamma_j}$$

The Kneser-Hecke operator  $K_k(p)$  also acts on this space via  $\Delta_m(K_k(p))$ .

**Theorem.** (Yoshida 1985) There are explicit constants c = c(m, k, p), d = d(m, k, p) such that

$$\delta_k(T_m^{(m-1)}(p^2)) = c \operatorname{id} + d\Delta_m(K_k(p)).$$

## Hecke operators as double cosets. (Codes.)

Let  $(R, A, \beta, Q)$  be a Type.

The associated extraspecial group

 $\mathcal{E}_m := (A^m \times A^m) \Join \mathbb{Q}/\mathbb{Z}, \text{ with multiplication} \\ (a, b, q)(a', b', q') = (a + a', b + b', q + q' + \beta(b', a))$ 

acts irreducibly on  $\mathbb{C}[A^m] = \langle x_v : v \in A^m \rangle_{\mathbb{C}}$  via

$$(a, b, q)x_v := \exp(2\pi i(q + \beta(v, a)))x_{v+b}$$

**Remark.** The associated Clifford-Weil group  $C_m \leq GL(\mathbb{C}[A^m])$  normalizes  $\mathcal{E}_m$ .

 $\mathcal{U}_j := \{(a, 0, 0) \mid a = (0^{m-j}, a_1, \dots, a_j) \in A^m\} \leq \mathcal{E}_m \text{ and } \mathcal{T}_j = \mathcal{C}_m p_{\mathcal{U}_j} \mathcal{C}_m$ where for  $U \leq \mathcal{E}_m$  the endomorphism

$$p_U := \frac{1}{|U|} \sum_{u \in U} u$$

denotes the orthogonal projection onto the fixed space of U. Note that  $p_U = 0$  if  $U \cap Z \neq \{(0,0,0)\}$  where  $Z = \{(0,0,q) \mid q \in \mathbb{Q}/\mathbb{Z}\} = Z(\mathcal{E}_m).$ 

**Theorem.** (N. 2006) If  $A = R = \mathbb{F}_q$  is a finite field, then

$$\mathcal{H}(\mathcal{C}_m) = \langle \mathcal{T}_j \mid 0 \le j \le m \rangle_{\mathbb{C}-algebra} = \mathbb{C}[\mathcal{T}_1]$$

is a commutative subalgebra of  $End(Inv(\mathcal{C}_m))$  consisting of selfadjoint linear operators acting on the subspace of degree N invariants via, say,  $\delta_N$ .

Then there are explicit constants c, d

(depending on q, the Type T, the genus m and the length N) such that

$$\delta_N(\mathcal{T}_1) = c \operatorname{id} + d\Delta_m(K_N(T)).$$