### Hecke Operators in Coding Theory

Elisabeth Nossek

**RWTH Aachen University** 

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### Codes

- $C \leq \mathbb{F}_q^N$  is called a linear code over  $\mathbb{F}_q$  of length N.
- $b: \mathbb{F}_q^N \times \mathbb{F}_q^N \mapsto \mathbb{F}_q: (x, y) \mapsto \sum_{i=1}^N x_i \overline{y_i}$  be a bilinear or hermitian form.
- $C^{\perp} := \{ v \in \mathbb{F}_q^N | b(c, v) = 0 \, \forall c \in C \}$  is the dual code of C.
- C is called selfdual if  $C^{\perp} = C$ , then dim(C) = N/2.

## **Kneser-Neighbors**

- C and D selfdual codes over 𝔽<sub>q</sub> with length N are called (1)-neighbors C ~ D iff dim(C/C ∩ D) = 1.
- neighboring graph  $\Gamma$ : vertices  $\mathcal{F} := \{C \leq \mathbb{F}_q^N | C = C^{\perp}\}$  and edges between neighbors.
- A adjacency matrix of Γ.

### Results about the neighboring graph

• **Theorem**:(Kneser)  $\Gamma$  is connected.

Determine all selfdual Codes in  $\mathcal{F}$  (or  $\mathcal{F}$ /equivalence) by going through the neighboring graph.



• **Theorem**:(Nebe)

Description of eigenvalues and eigenspaces of A.

### Higher Neighboring Relations

- C and D selfdual codes over 𝔽<sub>q</sub> with length N are called k-neighbors C ∼<sub>k</sub> D iff dim(C/C ∩ D) = k.
- k-neighboring graph  $\Gamma_k$ ,  $A_k$  adjacency matrix of  $\Gamma_k$ .
- Question: Are there polynomials  $p_k \in \mathbb{Q}[X]$  such that  $A_k = p_k(A)$ ?

### Adjacency Matrices

The powers of adjacency matrices "count" paths through the graph

 $(A^m)_{ij} = |\{\text{paths of length m from i-th to j-th vertex}\}|.$ 

# Isometry Group

- $G(\mathbb{F}_q^N, b)$  the group of isometries acts transitively on  $\mathcal{F} = \{C \leq \mathbb{F}_q^N | C = C^{\perp}\}$  (Witt's theorem).
- $G(\mathbb{F}_q^N, b)$  preserves neighboring relations.
- Hence  $|\{D \in \mathcal{F} | D \sim_k C\}|$  is independent of the choice of C.

# Example



# Decomposition of $A^k$

#### Theorem:

Let  $m_{kr}$  be the number of paths of length k between  $C \sim_r D$  in  $\Gamma_1$ 

$$A_1^k = \sum_{r=0}^k m_{kr} A_r$$

Recursive formula for  $m_{kr}$ :

$$m_{kr} = b_{r,r-1}m_{k-1,r-1} + b_{r,r}m_{k-1,r} + b_{r,r+1}m_{k-1,r+1}$$
  
where  $b_{ij} := |\{E \in \mathcal{F} | E \sim_j C \text{ and } E \sim_1 D\}$  for  $C \sim_i D$ .

### Example continued



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## Polynomial for $A_k$

### **Theorem:** Let $P = (p_{ij})_{0 \le i,j \le k} := M^{-1}$ with $M = (m_{ij})_{0 \le i,j \le k}$

$$A_k = p_k(A) := \sum_{j=0}^k p_{kj} A^j$$

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# $G(\mathbb{F}_q^N, b)$ -linearity

Let  $\mathcal{V} := \mathbb{C}^{\mathcal{F}} = \langle e_{\mathcal{C}} | \mathcal{C} \in \mathcal{F} \rangle_{\mathbb{C}}$  be a complex vector space:

V is a G(F<sup>N</sup><sub>q</sub>, b)-permutation module through the action of G(F<sup>N</sup><sub>q</sub>, b) on F

$$A_k: \mathcal{V} \to \mathcal{V}: e_C \mapsto \sum_{D \sim_k C} e_D$$

- $A_k$  is a  $G(\mathbb{F}_q^N, b)$ -linear endomorphism on  $\mathcal{V}$ .
- Theorem:  $\operatorname{End}_{\mathbb{C}G(\mathbb{F}_q^N,b)}(\mathcal{V}) = \mathbb{C}[A_1]$

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### **Hecke-Operators**

Let  $S_N$  act on  $\mathbb{F}_q^N$  by permuting the components.

- Let [C] be the orbit under the induced action on  $\mathcal F$
- $\overline{\mathcal{V}} := \langle [C] | C \in \mathcal{F} \rangle$
- Hecke operator:

$$T_k: \overline{\mathcal{V}} \to \overline{\mathcal{V}}: [C] \mapsto \sum_{D \sim_k C} [D]$$

$$A_k: \mathcal{V} \to \mathcal{V}: e_C \mapsto \sum_{D \sim_k C} e_D$$

• Corollary:  $p_k(T_1) = T_k$ 

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### Example for $T_k$

•  $\mathbb{F}_2$ , N = 8

• 2 equivalence classes:  $[e_8]$  and  $[i_2^4]$ .

•  $T_1 = \begin{bmatrix} 7 & 7 \\ 2 & 12 \end{bmatrix}$ .

$$T_{2} = \frac{1}{3}T_{1}^{2} - \frac{1}{3}T_{1} - \frac{14}{3}I_{2}$$
  
$$T_{3} = \frac{1}{27}T_{1}^{3} - \frac{4}{21}T_{1}^{2} - \frac{47}{21}T_{1} + 2I_{2}$$