

Actions, representations and various algebraic structures

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Actions and representations

An action of G on X is a map

$$A : X \times G \rightarrow X, \quad (x, g) \mapsto x \cdot g$$

A representation of G on X is a map

$$R : G \rightarrow X^X = \{f : X \rightarrow X\}$$

The two concepts are **the same**:

given A , set

$$R(g) := (x \mapsto A(x, g)) = (x \mapsto x \cdot g)$$

given R , set

$$A(x, g) := R(g)(x)$$

Group algebras — definition

Let \mathbb{F} be a field and G a finite group.

$\mathbb{F}G :=$ vector space with basis G , multiplication inherited from G and distributive law:

$$\left(\sum_{g \in G} \lambda_g \cdot g \right) \cdot \left(\sum_{\tilde{g} \in G} \mu_{\tilde{g}} \cdot \tilde{g} \right) = \sum_{g, \tilde{g} \in G} \lambda_g \cdot \mu_{\tilde{g}} \cdot (g\tilde{g})$$

for $\lambda_g, \mu_{\tilde{g}} \in \mathbb{F}$.

$\mathbb{F}G := \{f : G \rightarrow \mathbb{F}\}$ with pointwise addition and convolution product:

$$(f \cdot h)(g) := \sum_{\tilde{g} \in G} f(g \cdot \tilde{g}^{-1}) \cdot h(\tilde{g})$$

for $f, h : G \rightarrow \mathbb{F}$.

$\mathbb{F}G :=$ associative \mathbb{F} -algebra with generators G and relations $g \cdot \tilde{g} - (g\tilde{g}) = 0$ for $g, \tilde{g} \in G$.

Group algebras — properties

\mathbb{F} : field, G : group, $\mathbb{F}G$: group algebra, V : \mathbb{F} -vector space.

There is a **bijection** between

$$\{\varphi : G \rightarrow \text{GL}(V) \mid \varphi \text{ is a group homomorphism}\}$$

and

$$\{\psi : \mathbb{F}G \rightarrow \text{End}_{\mathbb{F}}(V) \mid \psi \text{ is an algebra homomorphism}\}$$

Given $\varphi : G \rightarrow \text{GL}(V)$, define

$$\psi \left(\sum_{g \in G} \lambda_g \cdot g \right) := \sum_{g \in G} \lambda_g \cdot \varphi(g)$$

(use **finite presentation**).

Given $\psi : \mathbb{F}G \rightarrow \text{End}_{\mathbb{F}}(V)$, simply restrict $\varphi := \psi|_G$, since

$$\mathbf{1}_V = \psi(\mathbf{1}_G) = \psi(g \cdot g^{-1}) = \psi(g) \cdot \psi(g^{-1}) \quad \text{for all } g \in G.$$

Modules

Definition (G -module or $\mathbb{F}G$ -module)

An \mathbb{F} -vector space V together with

- a group homomorphism $\varphi : G \rightarrow \text{GL}(V)$,
- or an algebra homomorphism $\psi : \mathbb{F}G \rightarrow \text{End}_{\mathbb{F}}(V)$

is called a G -module over \mathbb{F} or an $\mathbb{F}G$ -module.

This is nothing but

an \mathbb{F} -vector space with an \mathbb{F} -linear action for G .

This is nothing but

an \mathbb{F} -linear representation for G .

Kernels and faithfulness

Let $A : X \times G \rightarrow X$ be an **action**, or **equivalently**, let $R : G \rightarrow X^X$ be a **representation**.

Depending on the **types** of G and X , it might make sense to speak of the **kernel** of the representation R or not.

Definition (Faithful representation/action)

We call the **representation** R (or the **action** A) **faithful**, if its kernel $\ker R$ is trivial.

Note: If a G -module V over \mathbb{F} is faithful, **it does not necessarily follow** that the corresponding $\mathbb{F}G$ -module V is faithful!

Homomorphisms and isomorphisms

Let $A : X \times G \rightarrow X$ and $\tilde{A} : \tilde{X} \times G \rightarrow \tilde{X}$ be two actions.

Definition (G -homomorphism)

A **homomorphism** $\varphi : X \rightarrow \tilde{X}$ is called a G -homomorphism or G -equivariant, if

$$\varphi(x \cdot g) = \varphi(x) \cdot g \quad \text{for all } x \in X \text{ and all } g \in G.$$

Equivalently, this means

$$\varphi(A(x, g)) = \tilde{A}(\varphi(x), g) \quad \text{for all } x \in X \text{ and all } g \in G.$$

Equivalently, this means that this diagram commutes:

$$\begin{array}{ccc} X \times G & \xrightarrow{A} & X \\ \varphi \times \text{id}_G \downarrow & & \downarrow \varphi \\ \tilde{X} \times G & \xrightarrow{\tilde{A}} & \tilde{X} \end{array}$$

If φ has a G -equiv. inverse, it is called a G -isomorphism.

Subacts

Let G act on X , i.e. $A : X \times G \rightarrow X$.

Definition (G -invariant subset, Subact)

A subset $Y \subseteq X$ is called G -invariant, if

$$y \cdot g \in Y \quad \text{for all } y \in Y \text{ and all } g \in G.$$

The **restriction** $A|_{Y \times G}$ is then a map to Y and G acts on Y . If $Y \subseteq X$ is also a substructure of X , we call Y a **subact** (or **submodule** resp.).

Recall: A permutation representation was called **transitive** if it has no proper subacts.

Definition (Irreducible/simple module)

An $\mathbb{F}G$ -module M is called **irreducible** or **simple**, if it has no submodules except 0 and M itself.

Factor acts

Let G act on X , i.e. $A : X \times G \rightarrow X$.

Definition (G -invariant partition, factor act)

Let $X = \bigcup_{i \in I} Y_i$ be **partitioned** such that

$$\forall i \in I \text{ and } g \in G, \text{ we have } Y_i \cdot g \subseteq Y_j \text{ for some } j \in I.$$

We say that the **partition is G -invariant** and get an action on the set of parts $Y := \{Y_i \mid i \in I\}$:

$$Y_i * g := Y_j \quad \text{if} \quad Y_i \cdot g \subseteq Y_j.$$

Recall: We call a **permutation action primitive**, if it has no non-trivial factor acts.

Note: We usually want extra conditions to ensure that Y has the **same algebraic structure** as X and the new action is a **homomorphism** of such structures for all g .

Extensions and direct sums

This is only about modules!

Let

$$0 \longrightarrow W \xrightarrow{i} V \xrightarrow{\pi} U \cong V/W \longrightarrow 0$$

be a module V with a non-trivial submodule.

This sequence may or may not be **split**:

$$0 \longrightarrow W \xrightarrow{i} V \begin{array}{c} \xrightarrow{\pi} \\ \xleftarrow{r} \end{array} U \longrightarrow 0,$$

i.e. there is $r : U \rightarrow W$ with $\pi \circ r = \text{id}_U$.

If and only if it is **split**, the module V is isomorphic to the **direct sum**

$$V \cong W \oplus U.$$

Indecomposability and semisimplicity

Definition (Indecomposable module)

An $\mathbb{F}G$ -module V is called **indecomposable** if it is not isomorphic to a direct sum of two proper submodules. Otherwise it is called **decomposable**.

Lemma (Decomposable implies reducible)

A decomposable module is reducible.

Definition (Semisimple modules and algebras)

A module is called **semisimple**, if it is isomorphic to a direct sum of simple modules.

An \mathbb{F} -algebra \mathcal{A} is called **semisimple**, if every \mathcal{A} -module is semisimple.

Ordinary representation theory of groups

For a finite group, the group algebra $\mathbb{C}G$ is **semisimple**.

The ordinary representation theory of groups solves:

Problem (Classification of simple modules)

Classify the *isomorphism types* of **simple** $\mathbb{C}G$ -modules, *i.e. classify irreducible $\mathbb{C}G$ -modules up to isomorphism.*

Lemma (Characters)

Two representations

$$R_1 : G \rightarrow \mathrm{GL}(V) \quad \text{and} \quad R_2 : G \rightarrow \mathrm{GL}(W)$$

*afforded by two $\mathbb{C}G$ -modules V and W are **isomorphic**, if and only if their **characters** $\chi_1 = \mathrm{Tr} \circ R_1$ and $\chi_2 = \mathrm{Tr} \circ R_2$ are equal.*

*The two characters $\chi_i : G \rightarrow \mathbb{C}$ are **class functions**.*

Research problems in ordinary rep. theory

Already done:

- Character tables of **symmetric groups**.
- Character tables of **alternating groups**.
- The **ATLAS** (character tables of simple groups).
- Some **generic character tables**.

Still to do:

- Determine character tables for **more groups**.
- Determine more **generic tables** for whole families of groups.
- Devise **better algorithms** to compute tables.

Modular representation theory of groups

\mathbb{F} : field with $\text{char}(\mathbb{F}) \mid |G|$, then $\mathbb{F}G$ is **not semisimple**.

The modular rep. theory of groups strives to solve:

Problem (Classification of simple modules)

Classify the *isomorphism types* of **simple** $\mathbb{F}G$ -modules, *i.e. classify irreducible $\mathbb{F}G$ -modules up to isomorphism.*

Problem (Classification of indecomposable modules)

Classify the *isomorphism types* of **indecomposable** $\mathbb{F}G$ -modules.

Lemma (Brauer characters)

Two irreducible representations $R_1 : G \rightarrow \text{GL}(V)$ and $R_2 : G \rightarrow \text{GL}(W)$ *afforded* by two $\mathbb{F}G$ -modules V and W are **isomorphic**, if and only if their **Brauer characters** ψ_1 and ψ_2 are equal.

The two Brauer characters ψ_i take *values in \mathbb{C}* !

Research problems in modular rep. theory

Already done:

- Brauer tables of **some small** symmetric groups ($n \leq 18$).
- Brauer tables of **some small** alternating groups.
- **Modular ATLAS** (Brauer tables of simple groups). 1992 by Hiß, Jansen, Lux and Parker: groups up to page 100 in the ATLAS, now some more.

Still to do:

- Determine Brauer tables for **more groups**.
- Complete the **Modular ATLAS**.
- Classify **simple modules** of $\mathbb{F}S_n$.
- Compute the **2-modular Brauer table** of the **Monster**.
- Find an **algorithm** to compute a Brauer table???
- Classify **indecomposable $\mathbb{F}G$ -modules**???

Permutation groups

Problem (Permutation group algorithms)

Given $G := \langle g_1, \dots, g_k \in S_n \rangle \leq S_n$ *on a computer*.

Find *efficient* algorithms to compute with and in G :

- Test membership of $\pi \in S_n$ in G .
- Find the group order $|G|$.
- Decide whether $G = A_n$ or $G = S_n$ or none.
- Find orbits and blocks of primitivity.
- Find a presentation.
- Find the centre of G .
- ...

All of this is done and works well in **nearly linear time**:

runtime is bounded by $C \cdot n \cdot k \cdot \log^D(|G|)$.

Open questions for permutation groups

Still to do (in nearly linear time):

- Compute the **centraliser** $C_G(H)$ for some $H < S_n$.
- Compute the **derived subgroup** G' .
- Compute **intersections** of $G, H < S_n$.
- Compute **conjugacy classes** of permutation groups.
- Test $G, H < S_n$ for **conjugacy**.

Matrix and projective groups

Problem (Matrix group algorithms)

Given $G := \langle M_1, \dots, M_k \in \text{GL}(\mathbb{F}_q^n) \rangle \leq \text{GL}(\mathbb{F}_q^n)$ *on a computer*.

Ultimate goal: Answer similar questions as for permutation groups.

This is largely unsolved!

Problem (Projective group algorithms)

Given $G := \langle \bar{M}_1, \dots, \bar{M}_k \in \text{PGL}(n, q) \rangle \leq \text{PGL}(n, q)$ *on a computer*.

Ultimate goal: Answer similar questions as for permutation groups.

Constructive recognition

Problem (Constructive recognition)

Let \mathbb{F}_q be the field with q elements and

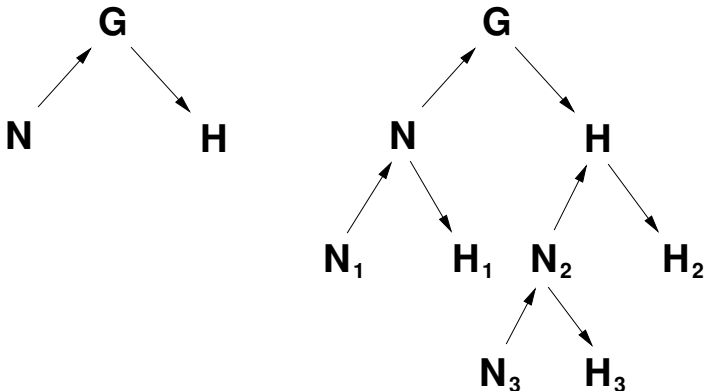
$$M_1, \dots, M_k \in \text{GL}(\mathbb{F}_q^n).$$

Find for $G := \langle M_1, \dots, M_k \rangle$:

- The group order $|G|$ and
- an algorithm that, given $M \in \text{GL}(\mathbb{F}_q^n)$,
 - *decides*, whether or not $M \in G$, and,
 - if so, expresses M *as word in the M_i* .
- The runtime should be bounded from above by a *polynomial in n , k and $\log q$* .
- A Monte Carlo Algorithmus is enough. (*Verification!*)

Recursion: composition trees

We get a tree:



Up arrows: **inclusions**

Down arrows: **homomorphisms**

Old idea, improvements are still being made

Enumerating large orbits

Orbit enumerations play an important role in

- modular representation theory,
- permutation group algorithms,
- matrix and projective group algorithms,
- combinatorics,
- finite geometry.

To get a feeling:

- To enumerate an orbit of 1140000 vectors in \mathbb{F}_2^{760} needs around **90 seconds**.
- To enumerate 95% of the same orbit with better **tricks** takes **1.1 seconds**.

Finding better ways to enumerate orbits is a current research topic.