



Families of correcting codes with ideal group algebra structure

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- C a q-ary Cauchy code of length q 1.
 - *C* is a left group code if and only one of the following conditions hold:
 - (a) C is permutation equivalent to $C_k(\alpha, f_m)$, with $L_{\alpha} = \mathbb{F}^*$ and $f_m(z) = z^m$.
 - (b) $2 \nmid q$ and C permutation equivalent to $C(\alpha, f_{m,m'})$, with $L_{\alpha} = \mathbb{F}^*$ and $f_{m,m'}(\xi^{2t+r}) = \xi^{2tm+rm'}$, $(\mathbb{F}^* = \langle \xi \rangle, t \in \mathbb{Z}, r \in \{0,1\}, 4m+k-1 \equiv 2m' \equiv 0 \mod (q-1)).$

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G group of order q - 1.
C is a left G-code if and only if either G is cyclic and condition (a) holds or G is dihedral and condition (b) holds.

 $C \subseteq \mathbb{FI}$ non-trivial affine-invariant code. Let a = a(C) and b = b(C). G group.

• C is a left G-code if and only if $G \simeq \mathcal{I}_{\alpha}$ for some map $\alpha : \mathcal{I} \to \mathcal{G}_{a,b}$ satisfying

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2 *C* is a *G*-code if and only if $G \simeq \mathcal{I}_{\alpha}$ for some map $\alpha : \mathcal{I} \to \operatorname{GL}(\mathbb{K}_{\mathbb{F}_{p^{a}}})$ satisfying (1) and such that the map $\beta : \mathbb{K} \times \mathbb{K} \to \mathbb{K}$ given by $\beta(x, y) = \alpha(x)^{-1}(y) - y$ is $\mathbb{F}_{p^{a}}$ -bilinear.