Matrix group recognition in GAP

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We can determine about G algorithmically (e.g.):

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Matrix groups in GAP

In standard GAP:

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[ <an immutable 56x56 matrix over GF2>,
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For “bigger” matrix groups both approaches do not work.

Constructive recognition — first formulation

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Find for $G := \langle M_1, \ldots, M_k \rangle$:

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If this problem is solved, we call

$\langle M_1, \ldots, M_k \rangle$ recognised constructively.

Complexity of algorithms

To measure the efficiency of an algorithm, we consider a class $\mathcal{P}$ of problems, that the algorithm can solve.

We assign to each $P \in \mathcal{P}$ its size $g(P)$,

and prove an upper bound for the runtime $L(P)$ of the algorithm for P:

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Example (Constructive matrix group recognition)

- Problem given by $M_1, \ldots, M_k \in \mathrm{GL}_n(\mathbb{F}_q)$.
- Size determined by n, k and $\log q$.
- Runtime should be $\leq$ a polynomial in n, k and $\log q$.

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Saving: 95% of runtime

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The discrete logarithm problem

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In practice q is small $\Rightarrow$ no problem.

We ignore both!

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- or at least “easier to recognise constructively”
- e.g. $H \leq S_m$ or $H \leq \mathrm{GL}_{n'}(\mathbb{F}_{q'})$ with $n' \log q' < n \log q$

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$\to$ Monte Carlo algorithm to compute N

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7. If $M \notin G$, then at least one step does not work.

Recursion: composition trees

We get a tree:

$$\begin{tikzcd} & \mathbf{G} \arrow[dr] & \\ \mathbf{N} \arrow[ur] & & \mathbf{H} \end{tikzcd}$$

Up arrows: inclusions

Down arrows: homomorphisms

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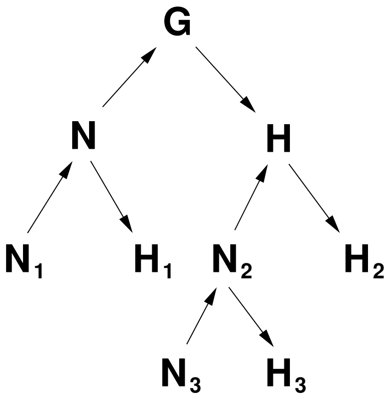

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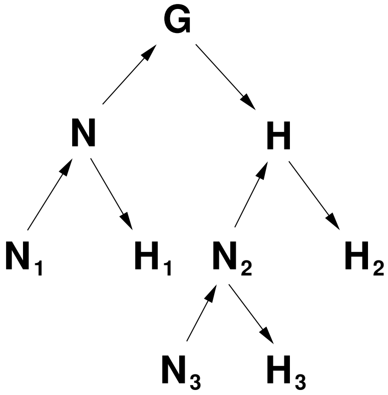

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Old idea, substantial improvements: Seress & N. 2006

Example: invariant subspace

Let $V = \mathbb{F}_q^n$, then G acts on V.
Let $W \leq V$ be an invariant subspace, i.e.:

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Choose basis $(w_1, \ldots, w_d)$ of W and extend to a basis

$$(w_1, \ldots, w_d, w_{d+1}, \ldots, w_n)$$

of V. After a base change the matrices in G look like this:

$$\left[\begin{array}{c|c} A & B \\ \hline \mathbf{0} & D \end{array}\right] \qquad \text{with } A \in \mathbb{F}_q^{d \times d}, B \in \mathbb{F}_q^{d \times (n-d)}, D \in \mathbb{F}_q^{(n-d) \times (n-d)}$$

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Together with a reduction additional information is gained!

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Theorem (Aschbacher, 1984)

Let $G \leq \mathrm{GL}_n(\mathbb{F}_q)$ and $Z := G \cap Z(\mathrm{GL}_n(\mathbb{F}_q))$ the subgroup of scalar matrices. Then G lies in at least one of the classes C1 to C8 or we have:

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This uses:

- the classification of finite simple groups
- information about their automorphism groups
- information about element orders
- information about conjugacy classes
- classifications of the irreducible representations
- information about the subgroup structure

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Usually this leads to Monte Carlo algorithms.

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Often this leads to efficient Las Vegas algorithms to find explicit isomorphisms.

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- Check the relations and thus prove the result.

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