

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Matrix group recognition

Max Neunhöffer

University of St Andrews

Aberdeen, 11.10.2007

Matrix groups ...

Let \mathbb{F}_q be the field with q elements and

$$\mathrm{GL}_n(\mathbb{F}_q) := \{M \in \mathbb{F}_q^{n \times n} \mid M \text{ invertible}\}$$

Introduction

Matrix groups

Constructive recognition

The problem

Complexity theory

Randomised algorithms

Constructive recognition

Troubles

Reduction

Homomorphisms

Computing the kernel

Recursion: composition
trees

Example: invariant
subspace

Finding reductions

Solution for leaves

Classifications

Recognition of the groups

Standard generators

Verification

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Matrix groups ...

Let \mathbb{F}_q be the field with q elements and

$$\mathrm{GL}_n(\mathbb{F}_q) := \{M \in \mathbb{F}_q^{n \times n} \mid M \text{ invertible}\}$$

Given: $M_1, \dots, M_k \in \mathrm{GL}_n(\mathbb{F}_q)$

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Matrix groups ...

Let \mathbb{F}_q be the field with q elements and

$$\mathrm{GL}_n(\mathbb{F}_q) := \{M \in \mathbb{F}_q^{n \times n} \mid M \text{ invertible}\}$$

Given: $M_1, \dots, M_k \in \mathrm{GL}_n(\mathbb{F}_q)$

Then the M_i generate a group $G \leq \mathrm{GL}_n(\mathbb{F}_q)$.

Matrix groups ...

Let \mathbb{F}_q be the field with q elements and

$$\mathrm{GL}_n(\mathbb{F}_q) := \{M \in \mathbb{F}_q^{n \times n} \mid M \text{ invertible}\}$$

Given: $M_1, \dots, M_k \in \mathrm{GL}_n(\mathbb{F}_q)$

Then the M_i generate a group $G \leq \mathrm{GL}_n(\mathbb{F}_q)$.

It is **finite**, we have $|\mathrm{GL}_n(\mathbb{F}_q)| = q^{n(n-1)/2} \prod_{i=1}^n (q^i - 1)$

Matrix groups ...

Let \mathbb{F}_q be the field with q elements and

$$\mathrm{GL}_n(\mathbb{F}_q) := \{M \in \mathbb{F}_q^{n \times n} \mid M \text{ invertible}\}$$

Given: $M_1, \dots, M_k \in \mathrm{GL}_n(\mathbb{F}_q)$

Then the M_i generate a group $G \leq \mathrm{GL}_n(\mathbb{F}_q)$.

It is **finite**, we have $|\mathrm{GL}_n(\mathbb{F}_q)| = q^{n(n-1)/2} \prod_{i=1}^n (q^i - 1)$

What do we want to determine about G ?

- The group order $|G|$

Matrix groups ...

Let \mathbb{F}_q be the field with q elements and

$$\mathrm{GL}_n(\mathbb{F}_q) := \{M \in \mathbb{F}_q^{n \times n} \mid M \text{ invertible}\}$$

Given: $M_1, \dots, M_k \in \mathrm{GL}_n(\mathbb{F}_q)$

Then the M_i generate a group $G \leq \mathrm{GL}_n(\mathbb{F}_q)$.

It is **finite**, we have $|\mathrm{GL}_n(\mathbb{F}_q)| = q^{n(n-1)/2} \prod_{i=1}^n (q^i - 1)$

What do we want to determine about G ?

- The group order $|G|$
- Membership test: Is $M \in \mathrm{GL}_n(\mathbb{F}_q)$ in G ?

Matrix groups ...

Let \mathbb{F}_q be the field with q elements and

$$\mathrm{GL}_n(\mathbb{F}_q) := \{M \in \mathbb{F}_q^{n \times n} \mid M \text{ invertible}\}$$

Given: $M_1, \dots, M_k \in \mathrm{GL}_n(\mathbb{F}_q)$

Then the M_i generate a group $G \leq \mathrm{GL}_n(\mathbb{F}_q)$.

It is **finite**, we have $|\mathrm{GL}_n(\mathbb{F}_q)| = q^{n(n-1)/2} \prod_{i=1}^n (q^i - 1)$

What do we want to determine about G ?

- The group order $|G|$
- Membership test: Is $M \in \mathrm{GL}_n(\mathbb{F}_q)$ in G ?
- Homomorphisms $\varphi : G \rightarrow H$?

Matrix groups ...

Let \mathbb{F}_q be the field with q elements and

$$\mathrm{GL}_n(\mathbb{F}_q) := \{M \in \mathbb{F}_q^{n \times n} \mid M \text{ invertible}\}$$

Given: $M_1, \dots, M_k \in \mathrm{GL}_n(\mathbb{F}_q)$

Then the M_i generate a group $G \leq \mathrm{GL}_n(\mathbb{F}_q)$.

It is **finite**, we have $|\mathrm{GL}_n(\mathbb{F}_q)| = q^{n(n-1)/2} \prod_{i=1}^n (q^i - 1)$

What do we want to determine about G ?

- The group order $|G|$
- Membership test: Is $M \in \mathrm{GL}_n(\mathbb{F}_q)$ in G ?
- Homomorphisms $\varphi : G \rightarrow H$?
- Kernels of homomorphisms? Is G simple?

Matrix groups ...

Let \mathbb{F}_q be the field with q elements and

$$\mathrm{GL}_n(\mathbb{F}_q) := \{M \in \mathbb{F}_q^{n \times n} \mid M \text{ invertible}\}$$

Given: $M_1, \dots, M_k \in \mathrm{GL}_n(\mathbb{F}_q)$

Then the M_i generate a group $G \leq \mathrm{GL}_n(\mathbb{F}_q)$.

It is **finite**, we have $|\mathrm{GL}_n(\mathbb{F}_q)| = q^{n(n-1)/2} \prod_{i=1}^n (q^i - 1)$

What do we want to determine about G ?

- The group order $|G|$
- Membership test: Is $M \in \mathrm{GL}_n(\mathbb{F}_q)$ in G ?
- Homomorphisms $\varphi : G \rightarrow H$?
- Kernels of homomorphisms? Is G simple?
- Comparison with known groups

Matrix groups ...

Let \mathbb{F}_q be the field with q elements and

$$\mathrm{GL}_n(\mathbb{F}_q) := \{M \in \mathbb{F}_q^{n \times n} \mid M \text{ invertible}\}$$

Given: $M_1, \dots, M_k \in \mathrm{GL}_n(\mathbb{F}_q)$

Then the M_i generate a group $G \leq \mathrm{GL}_n(\mathbb{F}_q)$.

It is **finite**, we have $|\mathrm{GL}_n(\mathbb{F}_q)| = q^{n(n-1)/2} \prod_{i=1}^n (q^i - 1)$

What do we want to determine about G ?

- The group order $|G|$
- Membership test: Is $M \in \mathrm{GL}_n(\mathbb{F}_q)$ in G ?
- Homomorphisms $\varphi : G \rightarrow H$?
- Kernels of homomorphisms? Is G simple?
- Comparison with known groups
- (Maximal) subgroups?

Matrix groups ...

Let \mathbb{F}_q be the field with q elements and

$$\mathrm{GL}_n(\mathbb{F}_q) := \{M \in \mathbb{F}_q^{n \times n} \mid M \text{ invertible}\}$$

Given: $M_1, \dots, M_k \in \mathrm{GL}_n(\mathbb{F}_q)$

Then the M_i generate a group $G \leq \mathrm{GL}_n(\mathbb{F}_q)$.

It is **finite**, we have $|\mathrm{GL}_n(\mathbb{F}_q)| = q^{n(n-1)/2} \prod_{i=1}^n (q^i - 1)$

What do we want to determine about G ?

- The group order $|G|$
- Membership test: Is $M \in \mathrm{GL}_n(\mathbb{F}_q)$ in G ?
- Homomorphisms $\varphi : G \rightarrow H$?
- Kernels of homomorphisms? Is G simple?
- Comparison with known groups
- (Maximal) subgroups?
- ...

Max Neunhöffer

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees

Example: invariant
subspace

Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Let \mathbb{F}_q be the field with q elements and

$$\mathrm{GL}_n(\mathbb{F}_q) := \{M \in \mathbb{F}_q^{n \times n} \mid M \text{ invertible}\}$$

Given: $M_1, \dots, M_k \in \mathrm{GL}_n(\mathbb{F}_q)$

Then the M_i generate a group $G \leq \mathrm{GL}_n(\mathbb{F}_q)$.

It is **finite**, we have $|\mathrm{GL}_n(\mathbb{F}_q)| = q^{n(n-1)/2} \prod_{i=1}^n (q^i - 1)$

Permutation groups and matrix groups

Max Neunhöffer

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees

Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Let $n \in \mathbb{N}$ and S_n be the **symmetric group**:

$$S_n = \{\pi : \{1, \dots, n\} \rightarrow \{1, \dots, n\} \mid \pi \text{ bijective}\}.$$

Let \mathbb{F}_q be the field with q elements and

$$\mathrm{GL}_n(\mathbb{F}_q) := \{M \in \mathbb{F}_q^{n \times n} \mid M \text{ invertible}\}$$

Given: $M_1, \dots, M_k \in \mathrm{GL}_n(\mathbb{F}_q)$

Then the M_i generate a group $G \leq \mathrm{GL}_n(\mathbb{F}_q)$.

It is **finite**, we have $|\mathrm{GL}_n(\mathbb{F}_q)| = q^{n(n-1)/2} \prod_{i=1}^n (q^i - 1)$

Permutation groups and matrix groups

Max Neunhoffer

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
treesExample: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Let $n \in \mathbb{N}$ and S_n be the **symmetric group**:

$$S_n = \{\pi : \{1, \dots, n\} \rightarrow \{1, \dots, n\} \mid \pi \text{ bijective}\}.$$

Given: $\pi_1, \dots, \pi_k \in S_n$

Let \mathbb{F}_q be the field with q elements and

$$\text{GL}_n(\mathbb{F}_q) := \{M \in \mathbb{F}_q^{n \times n} \mid M \text{ invertible}\}$$

Given: $M_1, \dots, M_k \in \text{GL}_n(\mathbb{F}_q)$

Then the M_i generate a group $G \leq \text{GL}_n(\mathbb{F}_q)$.

It is **finite**, we have $|\text{GL}_n(\mathbb{F}_q)| = q^{n(n-1)/2} \prod_{i=1}^n (q^i - 1)$

Permutation groups and matrix groups

Max Neunhöffer

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
treesExample: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Let $n \in \mathbb{N}$ and S_n be the **symmetric group**:

$$S_n = \{\pi : \{1, \dots, n\} \rightarrow \{1, \dots, n\} \mid \pi \text{ bijective}\}.$$

Given: $\pi_1, \dots, \pi_k \in S_n$

Then the π_i generate a group $G \leq S_n$.

Let \mathbb{F}_q be the field with q elements and

$$\text{GL}_n(\mathbb{F}_q) := \{M \in \mathbb{F}_q^{n \times n} \mid M \text{ invertible}\}$$

Given: $M_1, \dots, M_k \in \text{GL}_n(\mathbb{F}_q)$

Then the M_i generate a group $G \leq \text{GL}_n(\mathbb{F}_q)$.

It is **finite**, we have $|\text{GL}_n(\mathbb{F}_q)| = q^{n(n-1)/2} \prod_{i=1}^n (q^i - 1)$

Permutation groups and matrix groups

Max Neunhöffer

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
treesExample: invariant
subspace

Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Let $n \in \mathbb{N}$ and S_n be the **symmetric group**:

$$S_n = \{\pi : \{1, \dots, n\} \rightarrow \{1, \dots, n\} \mid \pi \text{ bijective}\}.$$

Given: $\pi_1, \dots, \pi_k \in S_n$

Then the π_i generate a group $G \leq S_n$.

It is **finite**, we have $|S_n| = n!$

Let \mathbb{F}_q be the field with q elements and

$$\text{GL}_n(\mathbb{F}_q) := \{M \in \mathbb{F}_q^{n \times n} \mid M \text{ invertible}\}$$

Given: $M_1, \dots, M_k \in \text{GL}_n(\mathbb{F}_q)$

Then the M_i generate a group $G \leq \text{GL}_n(\mathbb{F}_q)$.

It is **finite**, we have $|\text{GL}_n(\mathbb{F}_q)| = q^{n(n-1)/2} \prod_{i=1}^n (q^i - 1)$

Permutation groups

Let $n \in \mathbb{N}$ and S_n be the **symmetric group**:

$$S_n = \{\pi : \{1, \dots, n\} \rightarrow \{1, \dots, n\} \mid \pi \text{ bijective}\}.$$

Given: $\pi_1, \dots, \pi_k \in S_n$

Then the π_i generate a group $G \leq S_n$.

It is **finite**, we have $|S_n| = n!$.

We can determine about G algorithmically (e.g.):

- The group order $|G|$
- Membership test: Is $M \in S_n$ in G ?
- Homomorphisms $\varphi : G \rightarrow H$?
- Kernels of homomorphisms? Is G simple?
- Comparison with known groups
- (Maximal) subgroups?
- ...

Problem

Let \mathbb{F}_q be the field with q elements and

$$M_1, \dots, M_k \in \text{GL}_n(\mathbb{F}_q).$$

Find for $G := \langle M_1, \dots, M_k \rangle$:

- The group order $|G|$ and
- an algorithm that, given $M \in \text{GL}_n(\mathbb{F}_q)$,
 - **decides**, whether or not $M \in G$ and
 - if so, expresses M **as word in the M_i** .

Problem

Let \mathbb{F}_q be the field with q elements and

$$M_1, \dots, M_k \in \text{GL}_n(\mathbb{F}_q).$$

Find for $G := \langle M_1, \dots, M_k \rangle$:

- The group order $|G|$ and
- an algorithm that, given $M \in \text{GL}_n(\mathbb{F}_q)$,
 - **decides**, whether or not $M \in G$ and
 - if so, expresses M **as word in the M_i** .

If this problem is solved, we call

$\langle M_1, \dots, M_k \rangle$ recognised constructively.

Complexity of algorithms

To measure the **efficiency** of an algorithm, we consider a class \mathcal{P} of problems, that the algorithm can solve.

We assign to each $P \in \mathcal{P}$ its size $g(P)$,

and prove an upper bound for the runtime $L(P)$ of the algorithm for P :

$$L(P) \leq f(g(P))$$

for some function f .

Complexity of algorithms

To measure the **efficiency** of an algorithm, we consider a class \mathcal{P} of problems, that the algorithm can solve.

We assign to each $P \in \mathcal{P}$ its size $g(P)$,

and prove an upper bound for the runtime $L(P)$ of the algorithm for P :

$$L(P) \leq f(g(P))$$

for some function f .

The **growth rate of f** measures the **complexity**.

Complexity of algorithms

Max Neunhöffer

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

To measure the **efficiency** of an algorithm, we consider a class \mathcal{P} of problems, that the algorithm can solve.

We assign to each $P \in \mathcal{P}$ its size $g(P)$,

and prove an upper bound for the runtime $L(P)$ of the algorithm for P :

$$L(P) \leq f(g(P))$$

for some function f .

The **growth rate of f** measures the **complexity**.

Example (Constructive matrix group recognition)

- Problem given by $M_1, \dots, M_k \in \text{GL}_n(\mathbb{F}_q)$.
- Size determined by n , k and $\log q$.
- Runtime should be \leq a **polynomial** in n , k and $\log q$.

Max Neunhöffer

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Definition (Monte Carlo algorithms)

A Monte Carlo algorithm with error probability ϵ is an algorithm, that is **guaranteed** to terminate after a finite time, such that the **probability** that it returns a **wrong result** is at most ϵ .

Definition (Monte Carlo algorithms)

A Monte Carlo algorithm with error probability ϵ is an algorithm, that is **guaranteed** to terminate after a finite time, such that the **probability** that it returns a **wrong result** is at most ϵ .

Definition (Las Vegas algorithm)

A Las Vegas algorithm with error probability ϵ is an algorithm, that is **guaranteed** to terminate after a finite time, such that the **probability** that it **fails** is at most ϵ .

Randomised algorithms

Definition (Monte Carlo algorithms)

A Monte Carlo algorithm with error probability ϵ is an algorithm, that is **guaranteed** to terminate after a finite time, such that the **probability** that it returns a **wrong result** is at most ϵ .

Definition (Las Vegas algorithm)

A Las Vegas algorithm with error probability ϵ is an algorithm, that is **guaranteed** to terminate after a finite time, such that the **probability** that it **fails** is at most ϵ .

Example: Comp. of $|G| = 4\,089\,470\,473\,293\,004\,800$ for permutation group $G = \langle \pi_1, \pi_2 \rangle$ ($n = 137\,632$):

Randomised algorithms

Definition (Monte Carlo algorithms)

A Monte Carlo algorithm with error probability ϵ is an algorithm, that is **guaranteed** to terminate after a finite time, such that the **probability** that it returns a **wrong result** is at most ϵ .

Definition (Las Vegas algorithm)

A Las Vegas algorithm with error probability ϵ is an algorithm, that is **guaranteed** to terminate after a finite time, such that the **probability** that it **fails** is at most ϵ .

Example: Comp. of $|G| = 4\,089\,470\,473\,293\,004\,800$ for permutation group $G = \langle \pi_1, \pi_2 \rangle$ ($n = 137\,632$):
deterministic alg.: 112s

Definition (Monte Carlo algorithms)

A Monte Carlo algorithm with error probability ϵ is an algorithm, that is **guaranteed** to terminate after a finite time, such that the **probability** that it returns a **wrong result** is at most ϵ .

Definition (Las Vegas algorithm)

A Las Vegas algorithm with error probability ϵ is an algorithm, that is **guaranteed** to terminate after a finite time, such that the **probability** that it **fails** is at most ϵ .

Example: Comp. of $|G| = 4\,089\,470\,473\,293\,004\,800$ for permutation group $G = \langle \pi_1, \pi_2 \rangle$ ($n = 137\,632$):
deterministic alg.: 112s **Monte Carlo $\epsilon = 1\%$:** 6s
Saving: 95% of runtime

Constructive recognition

Problem

Let \mathbb{F}_q be the field with q elements and

$$M_1, \dots, M_k \in \text{GL}_n(\mathbb{F}_q).$$

Find for $G := \langle M_1, \dots, M_k \rangle$:

- The group order $|G|$ and
- an **algorithm** that, given $M \in \text{GL}_n(\mathbb{F}_q)$,
 - **decides**, whether or not $M \in G$, and,
 - if so, expresses M **as word in the M_i** .

Problem

Let \mathbb{F}_q be the field with q elements and

$$M_1, \dots, M_k \in \text{GL}_n(\mathbb{F}_q).$$

Find for $G := \langle M_1, \dots, M_k \rangle$:

- The group order $|G|$ and
- an **algorithm** that, given $M \in \text{GL}_n(\mathbb{F}_q)$,
 - **decides**, whether or not $M \in G$, and,
 - if so, expresses M **as word in the M_i** .
- The **runtime** should be bounded from above by a **polynomial in n , k and $\log q$** .

Problem

Let \mathbb{F}_q be the field with q elements and

$$M_1, \dots, M_k \in \text{GL}_n(\mathbb{F}_q).$$

Find for $G := \langle M_1, \dots, M_k \rangle$:

- The group order $|G|$ and
- an **algorithm** that, given $M \in \text{GL}_n(\mathbb{F}_q)$,
 - **decides**, whether or not $M \in G$, and,
 - if so, expresses M **as word in the M_i** .
- The **runtime** should be bounded from above by a **polynomial in n , k and $\log q$** .
- A **Monte Carlo Algorithmus** is enough.

Constructive recognition

Problem

Let \mathbb{F}_q be the field with q elements and

$$M_1, \dots, M_k \in \text{GL}_n(\mathbb{F}_q).$$

Find for $G := \langle M_1, \dots, M_k \rangle$:

- The group order $|G|$ and
- an **algorithm** that, given $M \in \text{GL}_n(\mathbb{F}_q)$,
 - **decides**, whether or not $M \in G$, and,
 - if so, expresses M **as word in the M_i** .
- The **runtime** should be bounded from above by a **polynomial in n , k and $\log q$** .
- A **Monte Carlo Algorithmus** is enough. **(Verification!)**

Problem

Let \mathbb{F}_q be the field with q elements and

$$M_1, \dots, M_k \in \text{GL}_n(\mathbb{F}_q).$$

Find for $G := \langle M_1, \dots, M_k \rangle$:

- The group order $|G|$ and
- an algorithm that, given $M \in \text{GL}_n(\mathbb{F}_q)$,
 - **decides**, whether or not $M \in G$, and,
 - if so, expresses M **as word in the M_i** .
- The **runtime** should be bounded from above by a **polynomial in n , k and $\log q$** .
- A Monte Carlo Algorithmus is enough. (**Verification!**)

If this problem is solved, we call

$\langle M_1, \dots, M_k \rangle$ **recognised constructively**.

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Troubles

The discrete logarithm problem

If $M_1 = [z] \in \mathbb{F}_q^{1 \times 1}$ with z a primitive root of \mathbb{F}_q . Then:

Given $0 \neq [x] \in \mathbb{F}_q^{1 \times 1}$, find $i \in \mathbb{N}$ such that $[x] = [z]^i$.

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Troubles

The discrete logarithm problem

If $M_1 = [z] \in \mathbb{F}_q^{1 \times 1}$ with z a primitive root of \mathbb{F}_q . Then:

Given $0 \neq [x] \in \mathbb{F}_q^{1 \times 1}$, find $i \in \mathbb{N}$ such that $[x] = [z]^i$.

There is no solution in polynomial time in $\log q$ known!

Troubles

The discrete logarithm problem

If $M_1 = [z] \in \mathbb{F}_q^{1 \times 1}$ with z a primitive root of \mathbb{F}_q . Then:

Given $0 \neq [x] \in \mathbb{F}_q^{1 \times 1}$, find $i \in \mathbb{N}$ such that $[x] = [z]^i$.

There is no solution in polynomial time in $\log q$ known!

Integer factorisation

Some methods need a factorisation of $q^i - 1$ for an $i \leq n$.

Troubles

The discrete logarithm problem

If $M_1 = [z] \in \mathbb{F}_q^{1 \times 1}$ with z a primitive root of \mathbb{F}_q . Then:

Given $0 \neq [x] \in \mathbb{F}_q^{1 \times 1}$, find $i \in \mathbb{N}$ such that $[x] = [z]^i$.

There is no solution in polynomial time in $\log q$ known!

Integer factorisation

Some methods need a factorisation of $q^i - 1$ for an $i \leq n$.

There is no solution in polynomial time in $\log q$ known!

Troubles

The discrete logarithm problem

If $M_1 = [z] \in \mathbb{F}_q^{1 \times 1}$ with z a primitive root of \mathbb{F}_q . Then:

Given $0 \neq [x] \in \mathbb{F}_q^{1 \times 1}$, find $i \in \mathbb{N}$ such that $[x] = [z]^i$.

There is no solution in polynomial time in $\log q$ known!

Integer factorisation

Some methods need a factorisation of $q^i - 1$ for an $i \leq n$.

There is no solution in polynomial time in $\log q$ known!

In practice q is small \Rightarrow no problem.

We ignore both!

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Let $G := \langle M_1, \dots, M_k \rangle \leq \text{GL}_n(\mathbb{F}_q)$.

What is a reduction?

Max Neunhöffer

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Let $G := \langle M_1, \dots, M_k \rangle \leq \text{GL}_n(\mathbb{F}_q)$.

A **reduction** is a group homomorphism

$$\begin{aligned} \varphi : G &\rightarrow H \\ M_i &\mapsto P_i \quad \text{for all } i \end{aligned}$$

with the following properties:

What is a reduction?

Max Neunhoffer

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Let $G := \langle M_1, \dots, M_k \rangle \leq \text{GL}_n(\mathbb{F}_q)$.

A **reduction** is a group homomorphism

$$\begin{aligned} \varphi : G &\rightarrow H \\ M_i &\mapsto P_i \quad \text{for all } i \end{aligned}$$

with the following properties:

- $\varphi(M)$ is **explicitly computable** for all $M \in G$

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

What is a reduction?

Let $G := \langle M_1, \dots, M_k \rangle \leq \text{GL}_n(\mathbb{F}_q)$.

A **reduction** is a group homomorphism

$$\begin{aligned} \varphi : G &\rightarrow H \\ M_i &\mapsto P_i \quad \text{for all } i \end{aligned}$$

with the following properties:

- $\varphi(M)$ is **explicitly computable** for all $M \in G$
- φ is **surjective**: $H = \langle P_1, \dots, P_k \rangle$

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

What is a reduction?

Let $G := \langle M_1, \dots, M_k \rangle \leq \text{GL}_n(\mathbb{F}_q)$.

A **reduction** is a group homomorphism

$$\begin{aligned} \varphi : G &\rightarrow H \\ M_i &\mapsto P_i \quad \text{for all } i \end{aligned}$$

with the following properties:

- $\varphi(M)$ is **explicitly computable** for all $M \in G$
- φ is **surjective**: $H = \langle P_1, \dots, P_k \rangle$
- H is in some sense “**smaller**”
- or at least “**easier to recognise constructively**”

Let $G := \langle M_1, \dots, M_k \rangle \leq \text{GL}_n(\mathbb{F}_q)$.

A **reduction** is a group homomorphism

$$\begin{aligned} \varphi : G &\rightarrow H \\ M_i &\mapsto P_i \quad \text{for all } i \end{aligned}$$

with the following properties:

- $\varphi(M)$ is **explicitly computable** for all $M \in G$
- φ is **surjective**: $H = \langle P_1, \dots, P_k \rangle$
- H is in some sense “**smaller**”
- or at least “**easier to recognise constructively**”
- e.g. $H \leq S_m$ or $H \leq \text{GL}_{n'}(\mathbb{F}_{q'})$ with $n' \log q' < n \log q$

Computing the kernel

Max Neunhoffer

Let $\varphi : G \rightarrow H$ be a reduction and assume that H is already recognised constructively.

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Computing the kernel

Let $\varphi : G \rightarrow H$ be a reduction and assume that H is already recognised constructively.

Then we can compute the kernel N of φ :

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Computing the kernel

Let $\varphi : G \rightarrow H$ be a reduction and assume that H is already recognised constructively.

Then we can compute the kernel N of φ :

1 Generate a (pseudo-) random element $M \in G$,

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Computing the kernel

Let $\varphi : G \rightarrow H$ be a reduction and assume that H is already recognised constructively.

Then we can compute the kernel N of φ :

- 1 Generate a (pseudo-) random element $M \in G$,
- 2 map it with φ onto $\varphi(M) \in H = \langle P_1, \dots, P_k \rangle$,

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Computing the kernel

Let $\varphi : G \rightarrow H$ be a reduction and assume that H is already recognised constructively.

Then we can compute the kernel N of φ :

- 1 Generate a (pseudo-) random element $M \in G$,
- 2 map it with φ onto $\varphi(M) \in H = \langle P_1, \dots, P_k \rangle$,
- 3 express $\varphi(M)$ as word in the P_i ,

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel

Recursion: composition
trees

Example: invariant
subspace

Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Computing the kernel

Let $\varphi : G \rightarrow H$ be a reduction and assume that H is already recognised constructively.

Then we can compute the kernel N of φ :

- 1 Generate a (pseudo-) random element $M \in G$,
- 2 map it with φ onto $\varphi(M) \in H = \langle P_1, \dots, P_k \rangle$,
- 3 express $\varphi(M)$ as word in the P_i ,
- 4 evaluate the same word in the M_i ,

Computing the kernel

Max Neunhöffer

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Let $\varphi : G \rightarrow H$ be a reduction and assume that H is already recognised constructively.

Then we can compute the kernel N of φ :

- 1 Generate a (pseudo-) random element $M \in G$,
- 2 map it with φ onto $\varphi(M) \in H = \langle P_1, \dots, P_k \rangle$,
- 3 express $\varphi(M)$ as word in the P_i ,
- 4 evaluate the same word in the M_i ,
- 5 get element $M' \in G$ with $M \cdot M'^{-1} \in N$.

Computing the kernel

Max Neunhöffer

Let $\varphi : G \rightarrow H$ be a reduction and assume that H is already recognised constructively.

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Then we can compute the kernel N of φ :

- 1 Generate a (pseudo-) random element $M \in G$,
- 2 map it with φ onto $\varphi(M) \in H = \langle P_1, \dots, P_k \rangle$,
- 3 express $\varphi(M)$ as word in the P_i ,
- 4 evaluate the same word in the M_i ,
- 5 get element $M' \in G$ with $M \cdot M'^{-1} \in N$.
- 6 If M is uniformly distributed in G
then $M \cdot M'^{-1}$ is uniformly distributed in N

Computing the kernel

Let $\varphi : G \rightarrow H$ be a reduction and assume that H is already recognised constructively.

Then we can compute the kernel N of φ :

- 1 Generate a (pseudo-) random element $M \in G$,
- 2 map it with φ onto $\varphi(M) \in H = \langle P_1, \dots, P_k \rangle$,
- 3 express $\varphi(M)$ as word in the P_i ,
- 4 evaluate the same word in the M_i ,
- 5 get element $M' \in G$ with $M \cdot M'^{-1} \in N$.
- 6 If M is uniformly distributed in G
then $M \cdot M'^{-1}$ is uniformly distributed in N
- 7 Repeat.

Computing the kernel

Max Neunhöffer

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Let $\varphi : G \rightarrow H$ be a reduction and assume that H is already recognised constructively.

Then we can compute the kernel N of φ :

- 1 Generate a (pseudo-) random element $M \in G$,
- 2 map it with φ onto $\varphi(M) \in H = \langle P_1, \dots, P_k \rangle$,
- 3 express $\varphi(M)$ as word in the P_i ,
- 4 evaluate the same word in the M_i ,
- 5 get element $M' \in G$ with $M \cdot M'^{-1} \in N$.
- 6 If M is uniformly distributed in G
then $M \cdot M'^{-1}$ is uniformly distributed in N
- 7 Repeat.

→ Monte Carlo algorithm to compute N

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel

Recursion: composition trees

Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Recognising image and kernel suffices

Let $\varphi : G \rightarrow H$ be a reduction and assume that **both** H **and** the kernel $N = \langle N_1, \dots, N_m \rangle$ of φ are already recognised constructively.

Recognising image and kernel suffices

Let $\varphi : G \rightarrow H$ be a reduction and assume that **both** H **and** the kernel $N = \langle N_1, \dots, N_m \rangle$ of φ are already recognised constructively.

Then we have recognised G constructively:

$$|G| = |H| \cdot |N|.$$

Recognising image and kernel suffices

Let $\varphi : G \rightarrow H$ be a reduction and assume that **both** H **and** the kernel $N = \langle N_1, \dots, N_m \rangle$ of φ are already recognised constructively.

Then we have recognised G constructively:

$$|G| = |H| \cdot |N|. \text{ And for } M \in \text{GL}_n(\mathbb{F}_q):$$

Recognising image and kernel suffices

Let $\varphi : G \rightarrow H$ be a reduction and assume that **both** H **and** the kernel $N = \langle N_1, \dots, N_m \rangle$ of φ are already recognised constructively.

Then we have recognised G constructively:

$$|G| = |H| \cdot |N|. \text{ And for } M \in \text{GL}_n(\mathbb{F}_q):$$

① **map** M with φ onto $\varphi(M) \in H = \langle P_1, \dots, P_k \rangle$,

Recognising image and kernel suffices

Let $\varphi : G \rightarrow H$ be a reduction and assume that **both** H **and** the kernel $N = \langle N_1, \dots, N_m \rangle$ of φ are already recognised constructively.

Then we have recognised G constructively:

$$|G| = |H| \cdot |N|. \text{ And for } M \in \text{GL}_n(\mathbb{F}_q):$$

- 1 map M with φ onto $\varphi(M) \in H = \langle P_1, \dots, P_k \rangle$,
- 2 express $\varphi(M)$ as word in the P_i ,

Recognising image and kernel suffices

Let $\varphi : G \rightarrow H$ be a reduction and assume that **both** H **and** the kernel $N = \langle N_1, \dots, N_m \rangle$ of φ are already recognised constructively.

Then we have recognised G constructively:

$$|G| = |H| \cdot |N|. \text{ And for } M \in \text{GL}_n(\mathbb{F}_q):$$

- 1 map M with φ onto $\varphi(M) \in H = \langle P_1, \dots, P_k \rangle$,
- 2 express $\varphi(M)$ as word in the P_i ,
- 3 evaluate the same word in the M_i ,

Recognising image and kernel suffices

Let $\varphi : G \rightarrow H$ be a reduction and assume that **both** H **and** the kernel $N = \langle N_1, \dots, N_m \rangle$ of φ are already recognised constructively.

Then we have recognised G constructively:

$$|G| = |H| \cdot |N|. \text{ And for } M \in \text{GL}_n(\mathbb{F}_q):$$

- 1 map M with φ onto $\varphi(M) \in H = \langle P_1, \dots, P_k \rangle$,
- 2 express $\varphi(M)$ as word in the P_i ,
- 3 evaluate the same word in the M_i ,
- 4 get element $M' \in G$ such that $M \cdot M'^{-1} \in N$,

Recognising image and kernel suffices

Let $\varphi : G \rightarrow H$ be a reduction and assume that **both** H **and** the kernel $N = \langle N_1, \dots, N_m \rangle$ of φ are already recognised constructively.

Then we have recognised G constructively:

$|G| = |H| \cdot |N|$. And for $M \in \text{GL}_n(\mathbb{F}_q)$:

- 1 map M with φ onto $\varphi(M) \in H = \langle P_1, \dots, P_k \rangle$,
- 2 express $\varphi(M)$ as word in the P_i ,
- 3 evaluate the same word in the M_i ,
- 4 get element $M' \in G$ such that $M \cdot M'^{-1} \in N$,
- 5 express $M \cdot M'^{-1}$ as word in the N_j ,

Recognising image and kernel suffices

Let $\varphi : G \rightarrow H$ be a reduction and assume that **both** H **and** the kernel $N = \langle N_1, \dots, N_m \rangle$ of φ are already recognised constructively.

Then we have recognised G constructively:

$$|G| = |H| \cdot |N|. \text{ And for } M \in \text{GL}_n(\mathbb{F}_q):$$

- 1 map M with φ onto $\varphi(M) \in H = \langle P_1, \dots, P_k \rangle$,
- 2 express $\varphi(M)$ as word in the P_i ,
- 3 evaluate the same word in the M_i ,
- 4 get element $M' \in G$ such that $M \cdot M'^{-1} \in N$,
- 5 express $M \cdot M'^{-1}$ as word in the N_j ,
- 6 get M as word in the M_i and N_j :
 $M' = \prod$ in the M_i , $M \cdot M'^{-1} = \prod$ in the N_j
 $\Rightarrow M = (\prod \text{ in the } N_j) \cdot (\prod \text{ in the } M_i).$

Recognising image and kernel suffices

Let $\varphi : G \rightarrow H$ be a reduction and assume that **both** H **and** the kernel $N = \langle N_1, \dots, N_m \rangle$ of φ are already recognised constructively.

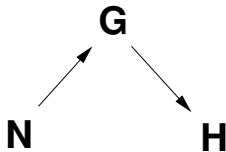
Then we have recognised G constructively:

$$|G| = |H| \cdot |N|. \text{ And for } M \in \text{GL}_n(\mathbb{F}_q):$$

- 1 map M with φ onto $\varphi(M) \in H = \langle P_1, \dots, P_k \rangle$,
- 2 express $\varphi(M)$ as word in the P_i ,
- 3 evaluate the same word in the M_i ,
- 4 get element $M' \in G$ such that $M \cdot M'^{-1} \in N$,
- 5 express $M \cdot M'^{-1}$ as word in the N_j ,
- 6 get M as word in the M_i and N_j :
 $M' = \prod$ in the M_i , $M \cdot M'^{-1} = \prod$ in the N_j
 $\Rightarrow M = (\prod \text{ in the } N_j) \cdot (\prod \text{ in the } M_i)$.
- 7 If $M \notin G$, then **at least** one step does not work.

Recursion: composition trees

We get a tree:



Up arrows: **inclusions**

Down arrows: **homomorphisms**

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel

Recursion: composition trees

Example: invariant
subspace
Finding reductions

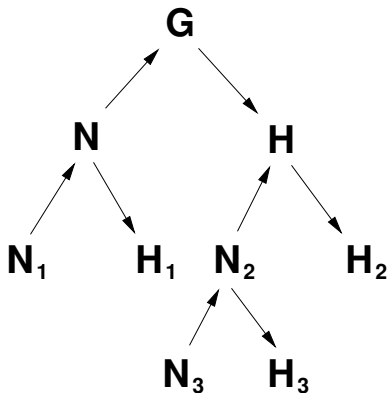
Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Recursion: composition trees

We get a tree:



Up arrows: inclusions

Down arrows: homomorphisms

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees

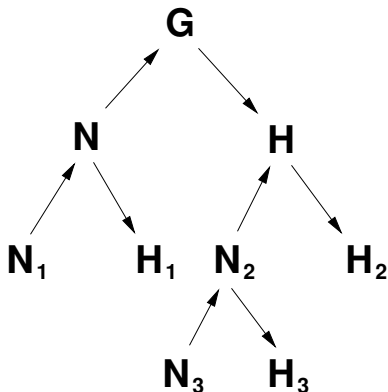
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

We get a tree:



Up arrows: **inclusions**

Down arrows: **homomorphisms**

Old idea, substantial improvements: Seress & N. 2006

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
**Example: invariant
subspace**
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Example: invariant subspace

Let $V = \mathbb{F}_q^n$, then G acts on V .

Let $W \leq V$ be an **invariant subspace**, i.e.:

$$MW = W \quad \text{for all } M \in G$$

Example: invariant subspace

Max Neunhoffer

Let $V = \mathbb{F}_q^n$, then G acts on V .Let $W \leq V$ be an invariant subspace, i.e.:

$$MW = W \quad \text{for all } M \in G$$

Choose basis (w_1, \dots, w_d) of W and extend to a basis

$$(w_1, \dots, w_d, w_{d+1}, \dots, w_n)$$

of V . After a base change the matrices in G look like this:

$$\left[\begin{array}{c|c} A & B \\ \hline \mathbf{0} & D \end{array} \right] \quad \text{with } A \in \mathbb{F}_q^{d \times d}, B \in \mathbb{F}_q^{d \times (n-d)}, D \in \mathbb{F}_q^{(n-d) \times (n-d)}$$

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Example: invariant subspace

Let $V = \mathbb{F}_q^n$, then G acts on V .

Let $W \leq V$ be an **invariant subspace**, i.e.:

$$MW = W \quad \text{for all } M \in G$$

Choose basis (w_1, \dots, w_d) of W and extend to a basis

$$(w_1, \dots, w_d, w_{d+1}, \dots, w_n)$$

of V . After a **base change** the matrices in G look like this:

$$\left[\begin{array}{c|c} A & B \\ \hline \mathbf{0} & D \end{array} \right] \quad \text{with } A \in \mathbb{F}_q^{d \times d}, B \in \mathbb{F}_q^{d \times (n-d)}, D \in \mathbb{F}_q^{(n-d) \times (n-d)}$$

and

$$G \rightarrow \text{GL}_{n-d}(\mathbb{F}_q), \left[\begin{array}{c|c} A & B \\ \hline \mathbf{0} & D \end{array} \right] \mapsto D$$

is a homomorphism of groups.

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
**Example: invariant
subspace**
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Example: invariant subspace

$$G \rightarrow \text{GL}_{n-d}(\mathbb{F}_q), \begin{bmatrix} A & B \\ \mathbf{0} & D \end{bmatrix} \mapsto D$$

is a homomorphism of groups, its kernel is

$$N := \left\{ \begin{bmatrix} A & B \\ \mathbf{0} & D \end{bmatrix} \in G \mid D = \mathbf{1} \right\}.$$

Example: invariant subspace

Max Neunhoffer

$$G \rightarrow \mathrm{GL}_{n-d}(\mathbb{F}_q), \begin{bmatrix} A & B \\ \mathbf{0} & D \end{bmatrix} \mapsto D$$

is a homomorphism of groups, its kernel is

$$N := \left\{ \begin{bmatrix} A & B \\ \mathbf{0} & D \end{bmatrix} \in G \mid D = \mathbf{1} \right\}.$$

The mapping

$$N \rightarrow \mathrm{GL}_d(\mathbb{F}_q), \begin{bmatrix} A & B \\ \mathbf{0} & \mathbf{1} \end{bmatrix} \mapsto A$$

also is a homomorphism of groups and has kernel

$$N_2 := \left\{ \begin{bmatrix} A & B \\ \mathbf{0} & D \end{bmatrix} \in G \mid A = D = \mathbf{1} \right\}.$$

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
treesExample: invariant
subspace

Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Example: invariant subspace

Max Neunhöffer

$$G \rightarrow \mathrm{GL}_{n-d}(\mathbb{F}_q), \left[\begin{array}{cc} A & B \\ \mathbf{0} & D \end{array} \right] \mapsto D$$

is a homomorphism of groups, its kernel is

$$N := \left\{ \left[\begin{array}{cc} A & B \\ \mathbf{0} & D \end{array} \right] \in G \mid D = \mathbf{1} \right\}.$$

The mapping

$$N \rightarrow \mathrm{GL}_d(\mathbb{F}_q), \left[\begin{array}{cc} A & B \\ \mathbf{0} & \mathbf{1} \end{array} \right] \mapsto A$$

also is a homomorphism of groups and has kernel

$$N_2 := \left\{ \left[\begin{array}{cc} A & B \\ \mathbf{0} & D \end{array} \right] \in G \mid A = D = \mathbf{1} \right\}.$$

This group is a p -group for $q = p^e$:

$$\left[\begin{array}{cc} \mathbf{1} & B \\ \mathbf{0} & \mathbf{1} \end{array} \right] \cdot \left[\begin{array}{cc} \mathbf{1} & B' \\ \mathbf{0} & \mathbf{1} \end{array} \right] = \left[\begin{array}{cc} \mathbf{1} & B + B' \\ \mathbf{0} & \mathbf{1} \end{array} \right]$$

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
treesExample: invariant
subspace

Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Example: invariant subspace

Max Neunhoffer

$$G \rightarrow \mathrm{GL}_{n-d}(\mathbb{F}_q), \begin{bmatrix} A & B \\ \mathbf{0} & D \end{bmatrix} \mapsto D$$

is a homomorphism of groups, its kernel is

$$N := \left\{ \begin{bmatrix} A & B \\ \mathbf{0} & D \end{bmatrix} \in G \mid D = \mathbf{1} \right\}.$$

The mapping

$$N \rightarrow \mathrm{GL}_d(\mathbb{F}_q), \begin{bmatrix} A & B \\ \mathbf{0} & \mathbf{1} \end{bmatrix} \mapsto A$$

also is a homomorphism of groups and has kernel

$$N_2 := \left\{ \begin{bmatrix} A & B \\ \mathbf{0} & D \end{bmatrix} \in G \mid A = D = \mathbf{1} \right\}.$$

This group is a p -group for $q = p^e$:

$$\begin{bmatrix} \mathbf{1} & B \\ \mathbf{0} & \mathbf{1} \end{bmatrix} \cdot \begin{bmatrix} \mathbf{1} & B' \\ \mathbf{0} & \mathbf{1} \end{bmatrix} = \begin{bmatrix} \mathbf{1} & B + B' \\ \mathbf{0} & \mathbf{1} \end{bmatrix}$$

Together with a reduction additional information is gained!

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
treesExample: invariant
subspace

Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Max Neunhöffer

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

How to find reductions?

Aschbacher has defined classes C1 to C8 of subgroups of $GL_n(\mathbb{F}_q)$.

How to find reductions?

Max Neunhöffer

Aschbacher has defined classes C1 to C8 of subgroups of $GL_n(\mathbb{F}_q)$.

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Theorem (Aschbacher, 1984)

Let $G \leq GL_n(\mathbb{F}_q)$ and $Z := G \cap Z(GL_n(\mathbb{F}_q))$ the subgroup of scalar matrices. Then G lies in *at least one* of the classes C1 to C8 *or* we have:

- $T \subseteq G/Z \subseteq \text{Aut}(T)$
for a non-abelian simple group T , *and*
- G acts absolutely irreducibly on $V = \mathbb{F}_q^n$.

How to find reductions?

Max Neunhöffer

Aschbacher has defined classes C1 to C8 of subgroups of $GL_n(\mathbb{F}_q)$.

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Theorem (Aschbacher, 1984)

Let $G \leq GL_n(\mathbb{F}_q)$ and $Z := G \cap Z(GL_n(\mathbb{F}_q))$ the subgroup of scalar matrices. Then G lies in *at least one* of the classes C1 to C8 *or* we have:

- $T \subseteq G/Z \subseteq \text{Aut}(T)$
for a non-abelian simple group T , *and*
- G acts absolutely irreducibly on $V = \mathbb{F}_q^n$.

(This last case is called C9.)

How to find reductions?

Max Neunhöffer

Aschbacher has defined classes C1 to C8 of subgroups of $GL_n(\mathbb{F}_q)$.

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition trees
Example: invariant subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Theorem (Aschbacher, 1984)

Let $G \leq GL_n(\mathbb{F}_q)$ and $Z := G \cap Z(GL_n(\mathbb{F}_q))$ the subgroup of scalar matrices. Then G lies in *at least one* of the classes C1 to C8 *or* we have:

- $T \subseteq G/Z \subseteq \text{Aut}(T)$
for a non-abelian simple group T , and
- G acts absolutely irreducibly on $V = \mathbb{F}_q^n$.

(This last case is called C9.)

Thus we can call in *heavy artillery*:

- the *classification of finite simple groups*

How to find reductions?

Aschbacher has defined classes C1 to C8 of subgroups of $GL_n(\mathbb{F}_q)$.

Theorem (Aschbacher, 1984)

Let $G \leq GL_n(\mathbb{F}_q)$ and $Z := G \cap Z(GL_n(\mathbb{F}_q))$ the subgroup of scalar matrices. Then G lies in *at least one* of the classes C1 to C8 *or* we have:

- $T \subseteq G/Z \subseteq \text{Aut}(T)$
for a non-abelian simple group T , *and*
- G acts absolutely irreducibly on $V = \mathbb{F}_q^n$.

(This last case is called C9.)

Thus we can call in *heavy artillery*:

- the *classification of finite simple groups*
- the *modular representation theory of simple groups*

Approach for leaves of the tree

If none of the algorithms for C1 to C8 has succeeded:

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Approach for leaves of the tree

If none of the algorithms for C1 to C8 has succeeded:

- 1 For “small” groups compute **direct isomorphism** onto a permutation group.

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Approach for leaves of the tree

If none of the algorithms for C1 to C8 has succeeded:

- 1 For “small” groups compute **direct isomorphism** onto a permutation group.
- 2 **Determine**, for which (simple) group $T \leq G/Z \leq \text{Aut}(T)$ holds.

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Approach for leaves of the tree

If none of the algorithms for C1 to C8 has succeeded:

- 1 For “small” groups compute **direct isomorphism** onto a permutation group.
- 2 **Determine**, for which (simple) group $T \leq G/Z \leq \text{Aut}(T)$ holds.
- 3 **Find** an explicit isomorphism onto a “standard copy” of an intermediate group S .

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Approach for leaves of the tree

If none of the algorithms for C1 to C8 has succeeded:

- 1 For “small” groups compute **direct isomorphism** onto a permutation group.
- 2 **Determine**, for which (simple) group $T \leq G/Z \leq \text{Aut}(T)$ holds.
- 3 **Find** an explicit isomorphism onto a “standard copy” of an intermediate group S .
- 4 Finally **use** information about S to **recognise** G **constructively**.

Approach for leaves of the tree

If none of the algorithms for C1 to C8 has succeeded:

- 1 For “**small**” groups compute **direct isomorphism** onto a permutation group.
- 2 **Determine**, for which (simple) group $T \leq G/Z \leq \text{Aut}(T)$ holds.
- 3 **Find** an explicit isomorphism onto a “standard copy” of an intermediate group S .
- 4 Finally **use** information about S to **recognise** G **constructively**.

This uses:

- the classification of **finite simple groups**
- information about their **automorphism groups**
- information about **element orders**
- information about **conjugacy classes**
- classifications of the **irreducible representations**
- information about the **subgroup structure**

Max Neunhöffer

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Methods for non-constructive recognition:

- Knowledge about representations narrows down the possibilities

Max Neunhöffer

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Methods for non-constructive recognition:

- Knowledge about representations narrows down the possibilities
- Statistics about orders of random elements

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Methods for non-constructive recognition:

- Knowledge about representations narrows down the possibilities
- Statistics about orders of random elements

Usually this leads to **Monte Carlo algorithms**.

Standard generators

In G we can only multiply, invert and compute orders.

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Standard generators

In G we can only multiply, invert and compute orders.
Suppose: $G \cong S$ with $T \leq S \leq \text{Aut}(T)$ and T simple.

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Standard generators

In G we can only multiply, invert and compute orders.

Suppose: $G \cong S$ with $T \leq S \leq \text{Aut}(T)$ and T simple.

Find a tuple $(s_1, \dots, s_r) \in S^r$ together with certain words p_1, \dots, p_m in the s_i , such that:

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Standard generators

In G we can only multiply, invert and compute orders.

Suppose: $G \cong S$ with $T \leq S \leq \text{Aut}(T)$ and T simple.

Find a tuple $(s_1, \dots, s_r) \in S^r$ together with certain words p_1, \dots, p_m in the s_i , such that:

- $S = \langle s_1, \dots, s_r \rangle$,

Standard generators

In G we can only multiply, invert and compute orders.

Suppose: $G \cong S$ with $T \leq S \leq \text{Aut}(T)$ and T simple.

Find a tuple $(s_1, \dots, s_r) \in S^r$ together with certain words p_1, \dots, p_m in the s_i , such that:

- $S = \langle s_1, \dots, s_r \rangle$,
- if $(s'_1, \dots, s'_r) \in S^r$ with
 - $|s_i| = |s'_i|$ for $1 \leq i \leq r$,
 - $|p_j| = |p'_j|$ for $1 \leq j \leq m$
(the p'_j are the same words in the s'_i),

Standard generators

In G we can only multiply, invert and compute orders.

Suppose: $G \cong S$ with $T \leq S \leq \text{Aut}(T)$ and T simple.

Find a tuple $(s_1, \dots, s_r) \in S^r$ together with certain words p_1, \dots, p_m in the s_i , such that:

- $S = \langle s_1, \dots, s_r \rangle$,
- if $(s'_1, \dots, s'_r) \in S^r$ with
 - $|s_i| = |s'_i|$ for $1 \leq i \leq r$,
 - $|p_j| = |p'_j|$ for $1 \leq j \leq m$
(the p'_j are the same words in the s'_i),

then $s_i \mapsto s'_i$ for $1 \leq i \leq r$ defines an automorphism of S .

Standard generators

In G we can only multiply, invert and compute orders.
Suppose: $G \cong S$ with $T \leq S \leq \text{Aut}(T)$ and T simple.

Find a tuple $(s_1, \dots, s_r) \in S^r$ together with certain words p_1, \dots, p_m in the s_i , such that:

- $S = \langle s_1, \dots, s_r \rangle$,
- if $(s'_1, \dots, s'_r) \in S^r$ with
 - $|s_i| = |s'_i|$ for $1 \leq i \leq r$,
 - $|p_j| = |p'_j|$ for $1 \leq j \leq m$
(the p'_j are the same words in the s'_i),

then $s_i \mapsto s'_i$ for $1 \leq i \leq r$ defines an automorphism of S .

Such elements are called “standard generators” of S .

Standard generators

In G we can only **multiply**, **invert** and **compute orders**.
Suppose: $G \cong S$ with $T \leq S \leq \text{Aut}(T)$ and T simple.

Find a tuple $(s_1, \dots, s_r) \in S^r$ together with certain words p_1, \dots, p_m in the s_i , such that:

- $S = \langle s_1, \dots, s_r \rangle$,
- if $(s'_1, \dots, s'_r) \in S^r$ with
 - $|s_i| = |s'_i|$ for $1 \leq i \leq r$,
 - $|p_j| = |p'_j|$ for $1 \leq j \leq m$
 (the p'_j are the same words in the s'_i),

then $s_i \mapsto s'_i$ for $1 \leq i \leq r$ defines an **automorphism** of S .

Such elements are called “**standard generators**” of S .

We **find** $G \cong S$ **explicitly** by finding a tuple (M_1, \dots, M_r) of standard generators in G .

Standard generators

Max Neunhöffer

In G we can only **multiply**, **invert** and **compute orders**.
Suppose: $G \cong S$ with $T \leq S \leq \text{Aut}(T)$ and T simple.

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
treesExample: invariant
subspace

Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Find a tuple $(s_1, \dots, s_r) \in S^r$ together with certain words p_1, \dots, p_m in the s_i , such that:

- $S = \langle s_1, \dots, s_r \rangle$,
- if $(s'_1, \dots, s'_r) \in S^r$ with
 - $|s_i| = |s'_i|$ for $1 \leq i \leq r$,
 - $|p_j| = |p'_j|$ for $1 \leq j \leq m$
 (the p'_j are the same words in the s'_i),

then $s_i \mapsto s'_i$ for $1 \leq i \leq r$ defines an **automorphism** of S .

Such elements are called “**standard generators**” of S .

We **find** $G \cong S$ **explicitly** by finding a tuple (M_1, \dots, M_r) of standard generators in G .

Often this leads to efficient **Las Vegas algorithms** to find **explicit isomorphisms**.

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Everywhere we used randomised methods:
Las Vegas and **Monte Carlo**.

⇒ **We have to check whether our result is correct!**

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Verification

Everywhere we used randomised methods:
Las Vegas and **Monte Carlo**.

⇒ **We have to check whether our result is correct!**

Idea:

- Find **(short) presentations** for the leaf-groups,

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Everywhere we used randomised methods:
Las Vegas and **Monte Carlo**.

⇒ **We have to check whether our result is correct!**

Idea:

- Find **(short) presentations** for the leaf-groups,
- put these together to one for the whole group.

Introduction

Matrix groups
Constructive recognition

The problem

Complexity theory
Randomised algorithms
Constructive recognition
Troubles

Reduction

Homomorphisms
Computing the kernel
Recursion: composition
trees
Example: invariant
subspace
Finding reductions

Solution for leaves

Classifications
Recognition of the groups
Standard generators

Verification

Everywhere we used randomised methods:
Las Vegas and **Monte Carlo**.

⇒ **We have to check whether our result is correct!**

Idea:

- Find **(short) presentations** for the leaf-groups,
- put these together to one for the whole group.
- Check the **relations** and thus prove the result.