

Actions and Reps

Max Neunhoffer

Actions and reps

Group algebras

Algebras

Modules

Faithfulness

Homomorphisms

Subacts

Factor acts

Extensions and
direct sums

Indecomposability

Problems

Ordinary rep. theory

Modular rep. theory

Permutation groups

Matrix and projective
groups

Orbits

Actions, representations and various algebraic structures

Max Neunhoffer



University of St Andrews

12 November 2008

Actions and representations

An action of G on X is a map

$$A : X \times G \rightarrow X, \quad (x, g) \mapsto x \cdot g$$

A representation of G on X is a map

$$R : G \rightarrow X^X = \{f : X \rightarrow X\}$$

The two concepts are **the same**:

given A , set

$$R(g) := (x \mapsto A(x, g)) = (x \mapsto x \cdot g)$$

given R , set

$$A(x, g) := R(g)(x)$$

Group algebras — definition

Let \mathbb{F} be a field and G a finite group.

$\mathbb{F}G :=$ vector space with basis G , multiplication inherited from G and distributive law:

$$\left(\sum_{g \in G} \lambda_g \cdot g \right) \cdot \left(\sum_{\tilde{g} \in G} \mu_{\tilde{g}} \cdot \tilde{g} \right) = \sum_{g, \tilde{g} \in G} \lambda_g \cdot \mu_{\tilde{g}} \cdot (g\tilde{g})$$

for $\lambda_g, \mu_{\tilde{g}} \in \mathbb{F}$.

$\mathbb{F}G := \{f : G \rightarrow \mathbb{F}\}$ with pointwise addition and convolution product:

$$(f \cdot h)(g) := \sum_{\tilde{g} \in G} f(g \cdot \tilde{g}^{-1}) \cdot h(\tilde{g})$$

for $f, h : G \rightarrow \mathbb{F}$.

$\mathbb{F}G :=$ associative \mathbb{F} -algebra with generators G and relations $g \cdot \tilde{g} - (g\tilde{g}) = 0$ for $g, \tilde{g} \in G$.

Group algebras — properties

\mathbb{F} : field, G : group, $\mathbb{F}G$: group algebra, V : \mathbb{F} -vector space.

There is a **bijection** between

$$\{\varphi : G \rightarrow \text{GL}(V) \mid \varphi \text{ is a } \mathbf{group\ homomorphism}\}$$

and

$$\{\psi : \mathbb{F}G \rightarrow \text{End}_{\mathbb{F}}(V) \mid \psi \text{ is an } \mathbf{algebra\ homomorphism}\}$$

Given $\varphi : G \rightarrow \text{GL}(V)$, define

$$\psi \left(\sum_{g \in G} \lambda_g \cdot g \right) := \sum_{g \in G} \lambda_g \cdot \varphi(g)$$

(use **finite presentation**).

Given $\psi : \mathbb{F}G \rightarrow \text{End}_{\mathbb{F}}(V)$, simply restrict $\varphi := \psi|_G$, since

$$\mathbf{1}_V = \psi(\mathbf{1}_G) = \psi(g \cdot g^{-1}) = \psi(g) \cdot \psi(g^{-1}) \quad \text{for all } g \in G.$$

Modules

Definition (G -module or $\mathbb{F}G$ -module)

An \mathbb{F} -vector space V together with

- a group homomorphism $\varphi : G \rightarrow \text{GL}(V)$,
- or an algebra homomorphism $\psi : \mathbb{F}G \rightarrow \text{End}_{\mathbb{F}}(V)$

is called a G -module over \mathbb{F} or an $\mathbb{F}G$ -module.

This is nothing but

an \mathbb{F} -vector space with an \mathbb{F} -linear action for G .

This is nothing but

an \mathbb{F} -linear representation for G .

Kernels and faithfulness

Let $A : X \times G \rightarrow X$ be an **action**, or **equivalently**, let $R : G \rightarrow X^X$ be a **representation**.

Depending on the **types** of G and X , it might make sense to speak of the **kernel** of the representation R or not.

Definition (Faithful representation/action)

We call the **representation** R (or the **action** A) **faithful**, if its kernel $\ker R$ is trivial.

Note: If a G -module V over \mathbb{F} is faithful, **it does not necessarily follow** that the corresponding $\mathbb{F}G$ -module V is faithful!

Homomorphisms and isomorphisms

Let $A : X \times G \rightarrow X$ and $\tilde{A} : \tilde{X} \times G \rightarrow \tilde{X}$ be two actions.

Definition (G -homomorphism)

A **homomorphism** $\varphi : X \rightarrow \tilde{X}$ is called a G -homomorphism or G -equivariant, if

$$\varphi(x \cdot g) = \varphi(x) \cdot g \quad \text{for all } x \in X \text{ and all } g \in G.$$

Equivalently, this means

$$\varphi(A(x, g)) = \tilde{A}(\varphi(x), g) \quad \text{for all } x \in X \text{ and all } g \in G.$$

Equivalently, this means that this diagram commutes:

$$\begin{array}{ccc} X \times G & \xrightarrow{A} & X \\ \varphi \times \text{id}_G \downarrow & & \downarrow \varphi \\ \tilde{X} \times G & \xrightarrow{\tilde{A}} & \tilde{X} \end{array}$$

If φ has a G -equiv. inverse, it is called a G -isomorphism.

Subacts

Let G act on X , i.e. $A : X \times G \rightarrow X$.

Definition (G -invariant subset, Subact)

A subset $Y \subseteq X$ is called G -invariant, if

$$y \cdot g \in Y \quad \text{for all } y \in Y \text{ and all } g \in G.$$

The **restriction** $A|_{Y \times G}$ is then a map to Y and G acts on Y . If $Y \subseteq X$ is also a substructure of X , we call Y a **subact** (or **submodule** resp.).

Recall: A permutation representation was called **transitive** if it has no proper subacts.

Definition (Irreducible/simple module)

An $\mathbb{F}G$ -module M is called **irreducible** or **simple**, if it has no submodules except 0 and M itself.

Factor acts

Let G act on X , i.e. $A : X \times G \rightarrow X$.

Definition (G -invariant partition, factor act)

Let $X = \bigcup_{i \in I} Y_i$ be **partitioned** such that

$$\forall i \in I \text{ and } g \in G, \text{ we have } Y_i \cdot g \subseteq Y_j \text{ for some } j \in I.$$

We say that the **partition is G -invariant** and get an action on the set of parts $Y := \{Y_i \mid i \in I\}$:

$$Y_i * g := Y_j \quad \text{if} \quad Y_i \cdot g \subseteq Y_j.$$

Recall: We call a **permutation action primitive**, if it has no non-trivial factor acts.

Note: We usually want extra conditions to ensure that Y has the **same algebraic structure** as X and the new action is a **homomorphism** of such structures for all g .

Extensions and direct sums

This is only about modules!

Let

$$0 \longrightarrow W \xrightarrow{i} V \xrightarrow{\pi} U \cong V/W \longrightarrow 0$$

be a module V with a non-trivial submodule.

This sequence may or may not be **split**:

$$0 \longrightarrow W \xrightarrow{i} V \begin{array}{c} \xrightarrow{\pi} \\ \xleftarrow{r} \end{array} U \longrightarrow 0,$$

i.e. there is $r : U \rightarrow W$ with $\pi \circ r = \text{id}_U$.

If and only if it is **split**, the module V is isomorphic to the **direct sum**

$$V \cong W \oplus U.$$

Indecomposability and semisimplicity

Definition (Indecomposable module)

An $\mathbb{F}G$ -module V is called **indecomposable** if it is not isomorphic to a direct sum of two proper submodules. Otherwise it is called **decomposable**.

Lemma (Decomposable implies reducible)

A decomposable module is reducible.

Definition (Semisimple modules and algebras)

A module is called **semisimple**, if it is isomorphic to a direct sum of simple modules.

An \mathbb{F} -algebra \mathcal{A} is called **semisimple**, if every \mathcal{A} -module is semisimple.

Ordinary representation theory of groups

For a finite group, the group algebra $\mathbb{C}G$ is **semisimple**.

The ordinary representation theory of groups solves:

Problem (Classification of simple modules)

Classify the *isomorphism types* of **simple** $\mathbb{C}G$ -modules, *i.e. classify irreducible $\mathbb{C}G$ -modules up to isomorphism.*

Lemma (Characters)

Two representations

$$R_1 : G \rightarrow \mathrm{GL}(V) \quad \text{and} \quad R_2 : G \rightarrow \mathrm{GL}(W)$$

*afforded by two $\mathbb{C}G$ -modules V and W are **isomorphic**, if and only if their **characters** $\chi_1 = \mathrm{Tr} \circ R_1$ and $\chi_2 = \mathrm{Tr} \circ R_2$ are equal.*

*The two characters $\chi_i : G \rightarrow \mathbb{C}$ are **class functions**.*

Research problems in ordinary rep. theory

Already done:

- Character tables of **symmetric groups**.
- Character tables of **alternating groups**.
- The **ATLAS** (character tables of simple groups).
- Some **generic character tables**.

Still to do:

- Determine character tables for **more groups**.
- Determine more **generic tables** for whole families of groups.
- Devise **better algorithms** to compute tables.

Modular representation theory of groups

\mathbb{F} : field with $\text{char}(\mathbb{F}) \mid |G|$, then $\mathbb{F}G$ is **not semisimple**.

The modular rep. theory of groups strives to solve:

Problem (Classification of simple modules)

Classify the *isomorphism types* of **simple** $\mathbb{F}G$ -modules, i.e. classify **irreducible** $\mathbb{F}G$ -modules up to isomorphism.

Problem (Classification of indecomposable modules)

Classify the *isomorphism types* of **indecomposable** $\mathbb{F}G$ -modules.

Lemma (Brauer characters)

Two irreducible representations $R_1 : G \rightarrow \text{GL}(V)$ and $R_2 : G \rightarrow \text{GL}(W)$ *afforded* by two $\mathbb{F}G$ -modules V and W are **isomorphic**, if and only if their **Brauer characters** ψ_1 and ψ_2 are equal.

The two Brauer characters ψ_i take *values* in \mathbb{C} !

Research problems in modular rep. theory

Already done:

- Brauer tables of **some small** symmetric groups ($n \leq 18$).
- Brauer tables of **some small** alternating groups.
- **Modular ATLAS** (Brauer tables of simple groups). 1992 by Hiß, Jansen, Lux and Parker: groups up to page 100 in the ATLAS, now some more.

Still to do:

- Determine Brauer tables for **more groups**.
- Complete the **Modular ATLAS**.
- Classify **simple modules** of $\mathbb{F}S_n$.
- Compute the **2-modular Brauer table** of the **Monster**.
- Find an **algorithm** to compute a Brauer table???
- Classify **indecomposable $\mathbb{F}G$ -modules**???

Permutation groups

Problem (Permutation group algorithms)

Given $G := \langle g_1, \dots, g_k \in S_n \rangle \leq S_n$ *on a computer*.

Find *efficient* algorithms to compute with and in G :

- Test membership of $\pi \in S_n$ in G .
- Find the group order $|G|$.
- Decide whether $G = A_n$ or $G = S_n$ or none.
- Find orbits and blocks of primitivity.
- Find a presentation.
- Find the centre of G .
- ...

All of this is done and works well in **nearly linear time**:

runtime is bounded by $C \cdot n \cdot k \cdot \log^D(|G|)$.

Open questions for permutation groups

Still to do (in nearly linear time):

- Compute the **centraliser** $C_G(H)$ for some $H < S_n$.
- Compute the **derived subgroup** G' .
- Compute **intersections** of $G, H < S_n$.
- Compute **conjugacy classes** of permutation groups.
- Test $G, H < S_n$ for **conjugacy**.

Matrix and projective groups

Problem (Matrix group algorithms)

Given $G := \langle M_1, \dots, M_k \in \text{GL}(\mathbb{F}_q^n) \rangle \leq \text{GL}(\mathbb{F}_q^n)$ *on a computer*.

Ultimate goal: Answer similar questions as for permutation groups.

This is largely unsolved!

Problem (Projective group algorithms)

Given $G := \langle \bar{M}_1, \dots, \bar{M}_k \in \text{PGL}(n, q) \rangle \leq \text{PGL}(n, q)$ *on a computer*.

Ultimate goal: Answer similar questions as for permutation groups.

Constructive recognition

Problem (Constructive recognition)

Let \mathbb{F}_q be the field with q elements and

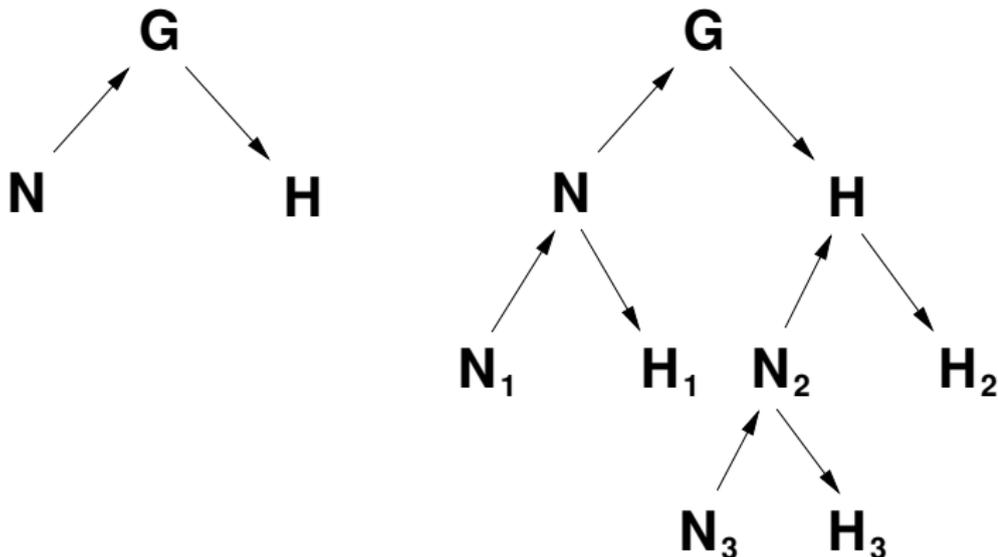
$$M_1, \dots, M_k \in \text{GL}(\mathbb{F}_q^n).$$

Find for $G := \langle M_1, \dots, M_k \rangle$:

- The group order $|G|$ and
- an algorithm that, given $M \in \text{GL}(\mathbb{F}_q^n)$,
 - *decides*, whether or not $M \in G$, and,
 - if so, expresses M *as word in the M_i* .
- The runtime should be bounded from above by a *polynomial in n , k and $\log q$* .
- A Monte Carlo Algorithmus is enough. (*Verification!*)

Recursion: composition trees

We get a tree:



Up arrows: **inclusions**

Down arrows: **homomorphisms**

Old idea, improvements are still being made

Enumerating large orbits

Orbit enumerations play an important role in

- modular representation theory,
- permutation group algorithms,
- matrix and projective group algorithms,
- combinatorics,
- finite geometry.

To get a feeling:

- To enumerate an orbit of 1140000 vectors in \mathbb{F}_2^{760} needs around **90 seconds**.
- To enumerate 95% of the same orbit with better **tricks** takes **1.1 seconds**.

Finding better ways to enumerate orbits is a current research topic.