The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

Enumerating orbits of Co_1 on $\mathbb{P}(\mathbb{F}_5^{24})$

Lehrstuhl D für Mathematik
RWTH Aachen

Perth 2006

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

All of this is joint work with:

- Robert A. Wilson
- Felix Noeske
- Jürgen Müller
- Frank Lübeck
- Christoph Köhler

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

Long orbits

Gunter Malle classified long orbits of quasi-simple groups:

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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Definition (Long orbit)

G: quasi-simple group, $\rho : G \to \mathrm{End}_{\mathbb{F}_q}(\mathbb{F}_q^d)$,
induces an action on the projective space $\mathbb{P}(\mathbb{F}_q^d)$

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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Recently he posed the following question:

Question

Does $2.Co_1$ have a long orbit in its action on $\mathbb{F}_5^{24}$?

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

The Action

$2.\mathrm{Co}_1 = \mathrm{Aut}(\Lambda)$ where Λ is the Leech lattice:

the unique 24-dimensional, even, unimodular lattice with no vectors of norm 2.

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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$2.Co_1 = \mathrm{Aut}(\Lambda)$ where Λ is the Leech lattice:

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Co_1 : one of the 26 sporadic simple groups.

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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$\Longrightarrow$ 24-dimensional integral representation of $2.Co_1$

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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We consider this representation mod 5.

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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We consider this representation mod 5.

$\longrightarrow$ Download two matrices in $\mathbb{F}_5^{24\times 24}$ from Rob's page.

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

The Size

$$|Co_1| = 4\,157\,776\,806\,543\,360\,000 \approx 4 \cdot 10^{18}$$

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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$$|\mathrm{Co}_1| = 4\,157\,776\,806\,543\,360\,000 \approx 4 \cdot 10^{18}$$

$$|\mathbb{P}(\mathbb{F}_5^{24})| = \frac{5^{24}-1}{5-1} = 14\,901\,161\,193\,847\,656 \approx 15 \cdot 10^{15}$$

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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Is there an orbit of length at least

$$\frac{5^{23}-1}{5-1} = 2\,980\,232\,238\,769\,531 \approx 3 \cdot 10^{15} \quad ?$$

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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Storing a field element in 4 Bits, we would need at least

1 387 778 Gigabytes $\approx$ 1.4 Petabytes

of memory to simply store all elements of such an orbit.

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

Standard orbit enumeration

Algorithm (Orbit enumeration)

Input: $G = \langle g_1, \ldots, g_r \rangle$ acting on X, $x \in X$
set $l := [x]$
for z in l:
    for g in $[g_1, \ldots, g_r]$:
        if zg in l:
            compute stabiliser element
        else:
            append zg to l
Output: l and generators for $\mathrm{Stab}_G(x)$

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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Output: I and generators for $\mathrm{Stab}_G(x)$

We need to

- store all points in memory,

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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We need to

- store all points in memory,
- look up points efficiently, and

Enumerating orbits of Co_1 on $\mathbb{P}(\mathbb{F}_5^{24})$

The Problem
The Question
The Action
The Size

Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits

The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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Output: l and generators for $\mathrm{Stab}_G(x)$

We need to

- store all points in memory,
- look up points efficiently, and
- compute xG and $\mathrm{Stab}_G(x)$ without knowing $|G|$.

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

Storing U-suborbits

$U < G$ a helper subgroup $\longrightarrow$ archive U-suborbits!

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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We want:

- given $x \in X$, store xU and compute $|xU|$

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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We want:

- given $x \in X$, store xU and compute $|xU|$
- given $z \in X$, decide whether z lies in a stored xU

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To this end, let $\bar{} : X \to Y$ be a homomorphism of U-sets:

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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To this end, let $\bar{}: X \to Y$ be a homomorphism of U-sets:

- enumerate Y completely

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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To this end, let $\bar{\ } : X \to Y$ be a homomorphism of U-sets:

- enumerate Y completely
- choose one element in each U-orbit of Y arbitrarily

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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To this end, let $\bar{\ } : X \to Y$ be a homomorphism of U-sets:

- enumerate Y completely
- choose one element in each U-orbit of Y arbitrarily
- call these U-minimal

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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- enumerate Y completely
- choose one element in each U-orbit of Y arbitrarily
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- for $y \in Y$, store a $u_y \in U$ such that yu_y is U-minimal

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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- choose one element in each U-orbit of Y arbitrarily
- call these U-minimal
- for $y \in Y$, store a $u_y \in U$ such that yu_y is U-minimal
- for U-minimal $y \in Y$, store generators of $\mathrm{Stab}_U(y)$

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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- for U-minimal $y \in Y$, store generators of $\mathrm{Stab}_U(y)$
- call $x \in X$ U-minimal, if $\bar{x} \in Y$ is U-minimal

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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- call these U-minimal
- for $y \in Y$, store a $u_y \in U$ such that yu_y is U-minimal
- for U-minimal $y \in Y$, store generators of $\mathrm{Stab}_U(y)$
- call $x \in X$ U-minimal, if $\bar{x} \in Y$ is U-minimal

Algorithm

Store xU by storing all U-minimal elements in xU.

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

Storing U-suborbits II

If $x \in X$ is U-minimal (i.e. $\bar{x} \in Y$ is U-minimal), then

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

Storing U-suborbits II

If $x \in X$ is U-minimal (i.e. $\bar{x} \in Y$ is U-minimal), then

$x\mathrm{Stab}_U(\bar{x})$ is the set of U-minimal elements in xU

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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Algorithm (Storing xU)

Input: $x \in X$
look up $u_{\bar{x}}$ and compute $z := xu_{\bar{x}}$
enumerate and store $z\mathrm{Stab}_U(\bar{z})$
find $\mathrm{Stab}_U(z) \leq \mathrm{Stab}_U(\bar{z})$ and thus $|zU| = |xU|$

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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enumerate and store $z\mathrm{Stab}_U(\bar{z})$
find $\mathrm{Stab}_U(z) \leq \mathrm{Stab}_U(\bar{z})$ and thus $|zU| = |xU|$

Algorithm (Looking up $z \in X$)

Input: $z \in X$, some stored xU
look up $u_{\bar{z}}$ and compute $w := zu_{\bar{z}}$
look up w in list of stored points
$z \in xU$ iff w already stored

Orbit by suborbits

Algorithm (Orbit by suborbits)

Input: $G = \langle g_1, \dots, g_r \rangle$ acting on X, $x \in X$

store xU and set $l := [x]$

repeat forever:

  for z in l:

    for g in $[g_1, \dots, g_r]$:

      if zgU already stored:

        compute stabiliser element

      else:

        store zgU

        append zg to l

  exit if orbit and stabiliser ready

Output: l, U-suborbits, generators for $\mathrm{Stab}_G(x)$

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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store xU and set $l := [x]$
repeat forever:
    for z in l:
        for g in $[g_1, \ldots, g_r]$:
            if zgU already stored:
                compute stabiliser element
            else:
                store zgU
                append zg to l
    exit if orbit and stabiliser ready
    for z in l:
        for u in generators of U:
            append zu to l
Output: l, U-suborbits, generators for $\mathrm{Stab}_G(x)$

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

Using two helper subgroups

$U < V < G$ two helper subgroups

$$\underbrace{\hat{} : X \to Z}_{\text{hom. of } V\text{-sets}} \quad , \quad \bar{} : Z \to Y$$

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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$\Longrightarrow$ can archive U-suborbits in Z and X!

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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Preparations:

- enumerate Z completely by U-orbits

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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$U < V < G$ two helper subgroups

$$\underbrace{\hat{}: X \to Z}_{\text{hom. of } V\text{-sets}} \quad , \quad \underbrace{\bar{}: Z \to Y \quad , \quad \bar{}: X \to Z \to Y}_{\text{hom. of } U\text{-sets}}$$

$\Longrightarrow$ can archive U-suborbits in Z and X!

Preparations:

- enumerate Z completely by U-orbits
- choose one U-minimal point in each V-orbit of Z, call it V-minimal

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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- for every U-minimal point in $z \in Z$ store:
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  - nothing if the V-minimal point of zV lies in zU

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

V-minimalising

Algorithm (V-minimalisation)

Input: $x \in X$
lookup $u \in U$ such that $w := xu$ is U-minimal
($\Longrightarrow \hat{w} \in V$ is U-minimal)
if $\hat{w}U$ does not contain the V-minimal point of $\hat{w}V$:
    look up $t \in \mathcal{T}$ such that $\hat{w}tU$ contains it
    set $w := wt$
    look up $u' \in U$ such that wu' is U-minimal
    set $w := wu'$
else:
    set $t := 1$ and $u' := 1$

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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else:

  set $t := 1$ and $u' := 1$

(now w is U-minimal and

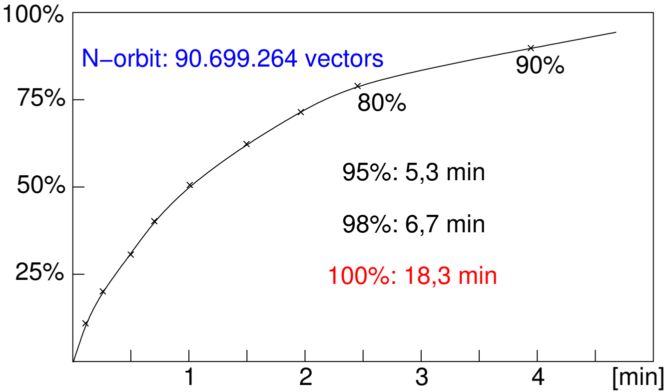
  $\hat{w}U$ contains the V-minimal point of $\hat{w}V$)

unless $\hat{w}$ is the V-minimal point:

  find $s \in \mathrm{Stab}_U(\bar{w})$ with $\hat{w}s$ V-minimal

Output: $v_x := utu's \in V$ such that xv_x is V-minimal

Evolution of an orbit enumeration



This is a typical time evolution for orbit enumerations!

Enumerating orbits of Co_1 on $\mathbb{P}(\mathbb{F}_5^{24})$

The Problem
The Question
The Action
The Size

Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits

The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

A half is enough!

Assume we

- know $|G|$,
- already have enumerated some part of xG, and
- already know some $S < \mathrm{Stab}_G(x)$ and $|S|$.

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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Assume we

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Then:

$$2 \cdot \text{Size(enumerated part)} \cdot |S| \geq |G|$$

if and only if

- S already is the full stabiliser $\mathrm{Stab}_G(x)$ and
- we already have enumerated at least half of $|xG|$

because if $S < \mathrm{Stab}_G(x)$ then the index is at least 2.

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

Finding homomorphisms

Let G act linearly on a F-vectorspace M:

$$\rho : G \to \mathrm{End}_F(M)$$

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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$N < M$ a G-invariant subspace,
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Then the following diagram commutes for all $g \in G$:

$$\begin{tikzcd} M \arrow[r, "\cdot g"] \arrow[d, "\pi"'] & M \arrow[d, "\pi"] \\ M/N \arrow[r, "\cdot g"] & M/N \end{tikzcd}$$

with the induced action on M/N.

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The same holds for the projective action, if we replace

- M by $\mathbb{P}(M)$ and
- $\mathbb{P}(M/N)$ by $\mathbb{P}(M/N) \cup \{0\}$.

Finding orbits

We can now enumerate halves of orbits.

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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We can now enumerate halves of orbits.

But how do we avoid enumerating them more than once?

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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Assume we "know" a half of xG, then for some $w \in X$ we can still check, whether $w \in xG$:

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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Algorithm (Membership test in half-orbit)

Input: $w \in X$ and at least a half of xG.
for 100 random elements $g \in G$:
    if wg in half of xG:
        return True
return False

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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Find bigger orbits by random search.

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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Find bigger orbits by random search.

But how to find small orbits?

Finding the small orbits

Short orbits have big stabilisers.

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

Finding the small orbits

Short orbits have big stabilisers.

Guess stabilisers:

- use maximal subgroups ($\to$ Rob's WWW Atlas)
- find invariant subspaces ($\to$ MEATAXE)

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

Finding the small orbits

Short orbits have big stabilisers.

Guess stabilisers:

- use maximal subgroups ($\rightarrow$ Rob's WWW Atlas)
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Guess elements of stabilisers:

- use conjugacy class reps. ($\rightarrow$ Rob's WWW Atlas)
- try vectors in eigenspaces

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

Finding all orbits

Build up a database of halves of pairwise disjoint orbits.

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

Finding all orbits

Build up a database of halves of pairwise disjoint orbits.

Produce representative candidates for the small orbits.

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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Build up a database of halves of pairwise disjoint orbits.

Produce representative candidates for the small orbits.

Produce random representatives for the big orbits.

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

Finding all orbits

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For all vectors: Test if they are in a known orbit half.
If not, enumerate half of new orbit.

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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For all vectors: Test if they are in a known orbit half.
If not, enumerate half of new orbit.

Do this until the sum of the orbit lengths is the total number of points.

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

Verification of Disjointness

How can we prove that two orbits are different knowing only half of them?

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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How can we prove that two orbits are different knowing only half of them?

Solution: Enumerate 51% of the orbits!

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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How can we prove that two orbits are different knowing only half of them?

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Lemma (Disjointness)

Two subsets of xG of size $> |xG|/2$ intersect nontrivially.

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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Lemma (Disjointness)

Two subsets of xG of size $> |xG|/2$ intersect nontrivially.

Algorithm (Disjointness proof)

Input: $M \subseteq xG$ with $2 \cdot |M| > |xG|$ and
$M' \subseteq x'G$ with $2 \cdot |M'| > |x'G|$
assume both M and M' are unions of V-sets

Check whether a V-orbit rep. of M is in M' or not.

The Result

The orbit lengths of the 48 orbits of Co_1 on $\mathbb{P}(\mathbb{F}_5^{24})$ are:

98 280	636 539 904 000	103 119 464 448 000
8 386 560	1 080 188 928 000	180 459 062 784 000
199 017 000	1 611 241 632 000	180 459 062 784 000
226 044 000	4 687 248 384 000	193 348 995 840 000
2 314 690 560	4 687 248 384 000	262 485 909 504 000
4 577 391 000	9 374 496 768 000	300 765 104 640 000
4 629 381 120	12 889 933 056 000	524 971 819 008 000
16 982 784 000	12 889 933 056 000	721 836 251 136 000
46 872 483 840	17 186 577 408 000	773 395 983 360 000
67 135 068 000	17 823 117 312 000	824 955 715 584 000
93 744 967 680	21 873 825 792 000	824 955 715 584 000
318 269 952 000	21 873 825 792 000	1 203 060 418 560 000
402 810 408 000	32 998 228 623 360	1 443 672 502 272 000
407 586 816 000	51 559 732 224 000	1 924 896 669 696 000
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Long limit: 2 980 232 238 769 531

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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$\Longrightarrow$ no long orbit!

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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199 017 000	1 611 241 632 000	180 459 062 784 000
226 044 000	4 687 248 384 000	193 348 995 840 000
2 314 690 560	4 687 248 384 000	262 485 909 504 000
4 577 391 000	9 374 496 768 000	300 765 104 640 000
4 629 381 120	12 889 933 056 000	524 971 819 008 000
16 982 784 000	12 889 933 056 000	721 836 251 136 000
46 872 483 840	17 186 577 408 000	773 395 983 360 000
67 135 068 000	17 823 117 312 000	824 955 715 584 000
93 744 967 680	21 873 825 792 000	824 955 715 584 000
318 269 952 000	21 873 825 792 000	1 203 060 418 560 000
402 810 408 000	32 998 228 623 360	1 443 672 502 272 000
407 586 816 000	51 559 732 224 000	1 924 896 669 696 000
407 586 816 000	69 296 280 109 056	2 165 508 753 408 000
563 934 571 200	103 119 464 448 000	2 887 345 004 544 000

Long limit: 2 980 232 238 769 531

$\Longrightarrow$ no long orbit!

Total: 14 901 161 193 847 656

The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
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The Result
Memory and Runtime Data

Memory and Runtime Data

In the end, the calculation

- used three helper subgroups: $U_1 < U_2 < U_3 < \mathrm{Co}_1$
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The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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The Problem
The Question
The Action
The Size
Enumerating large Orbits
Standard orbit enumeration
Using one helper subgroup
Orbit by suborbits
Using two helper subgroups
Halves of orbits
The Solution
Finding homomorphisms
Finding (small) orbits
Finding all orbits
Verification of Disjointness
The Result
Memory and Runtime Data

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