

Automorphism Groups of Hyperbolic Lattices

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1 Introduction

2 Preliminaries

3 Hyperbolic Lattices

4 Example

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Notation

- *Lattice*: a pair of free \mathbb{Z} -module L of rank n and a symmetric, non-degenerate bilinear form Φ on $\mathcal{V} = L \otimes_{\mathbb{Z}} \mathbb{R}$.
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- *Automorphism group*:

$$\text{Aut}(L) = \{g \in \text{GL}(\mathcal{V}) \mid Lg = L \text{ and } \Phi(xg, yg) = \Phi(x, y) \\ \text{for all } x, y \in L\}$$

$$\text{Aut}_{\mathbb{Z}}(A) = \{g \in \text{GL}_n(\mathbb{Z}) \mid gAg^{tr} = A\}$$

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For a given Gram matrix A of a lattice (L, Φ) determine a generating system for $\text{Aut}_{\mathbb{Z}}(A)$.

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- R.E. Borcherds, J.H. Conway ... (\sim 1988): results for special lattices, e.g. $(\mathbb{Z}^n, \text{diag}(-1, 1, \dots, 1))$ for small n .

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\mathcal{V} real vector space of dimension n , σ positive definite bilinear form on \mathcal{V} .
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(DC.1) $\sigma(x, y) > 0$ for all $x \in \mathcal{V}_1^{>0}, y \in \mathcal{V}_2^{>0}$.

(DC.2) For any $x \in \mathcal{V} \setminus \mathcal{V}_1^{>0}$ there is a $y \in \mathcal{V}_2^{>0} \setminus \{0\}$ such that
 $\sigma(x, y) \leq 0$. ($\mathcal{V}_i^{\geq 0} := \overline{\mathcal{V}_i^{>0}}$).

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(vi) $P_D = \{x \in \mathcal{V}_1^{\geq 0} \mid x \text{ } D\text{-perfect, } \mu_D(x) = 1\}$

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Proposition

$x \in P_D$, $r \notin \mathcal{V}_1^{\geq 0}$ direction of x .

$$\Rightarrow y = x + \rho r \in P_D \text{ for a } \rho > 0.$$

$x + \rho r$ is the *neighbour* of x in direction r . x, y are called *contiguous*.

D -Voronoi Graph

Γ_D graph with

$$V(\Gamma_D) = P_D$$

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- Γ_D is connected, i.e. the D -Voronoi-domains cover $\mathcal{V}_2^{>0}$ (exact face-to-face tessellation).

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- Γ_D is connected, i.e. the D -Voronoi-domains cover $\mathcal{V}_2^{>0}$ (exact face-to-face tessellation).
- Γ_D is locally finite, i.e. every D -perfect point has only finitely many neighbours.

Lemma

$\Omega \leq \mathrm{GL}(\mathcal{V})$ discrete subgroup acting properly discontinuously on $\mathcal{V}_1^{>0}$,
 $D\Omega^{ad} \subseteq D$. Then Ω acts on Γ_D as well.

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Theorem (Opgenorth, 2001)

Let $D \subset \mathcal{V}_2^{\geq 0} \setminus \{0\}$ admissible and discrete in \mathcal{V}_2 and $\Omega \leq \mathrm{GL}(\mathcal{V})$ as above and let the residue class graph Γ_D/Ω be finite. Further let $X := \{x_1, \dots, x_\ell\}$ a transversal of D -perfect points, which span a subtree T of Γ_D and let T_1 denote the neighbourhood of T in Γ_D . For all $y \in T_1$ choose an $\omega_y \in \Omega$, such that $\omega_y^{-1}(y) \in X$. Then

$$\Omega = \langle \omega_y, \mathrm{Stab}_\Omega x_i \mid y \in T_1, i = 1 \dots \ell \rangle.$$

In particular, Ω is finitely generated.

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Fix $A \in \mathbb{Z}_{hyp}^{n \times n}$ and $T \in \mathrm{GL}_n(\mathbb{R})$ s.t. $TAT^{tr} = \mathrm{diag}(-1, 1, \dots, 1)$
(Sylvester's Law of Inertia).

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- Dual cones w.r.t. standard scalar product:

$$\mathcal{V}_1^{>0} := \{x \in \mathbb{R}^n \mid xAx^{tr} < 0 \text{ and } xy_1^{tr} > 0\}$$

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- $D = \mathbb{Z}^n \cap \mathcal{V}_2^{\geq 0} \setminus \{0\}$
- $\Omega = \mathrm{Stab}_{\mathrm{Aut}_{\mathbb{Z}}(A)}(\mathcal{V}_1^{>0})$ (*reduced automorphism group*)

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- M_ε is convex
- M_ε is not a subset of a maximal convex lattice-free set (classified by Lovász, 1989).
 $\Rightarrow M_\varepsilon$ contains integral point.



D -minimal Vectors

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- continue automorphisms of $L(x)$ to rational automorphisms of A and find integral ones.
- connecting elements analogously computable

Theorem (M., 2013)

Let $A \in \mathbb{Z}_{hyp}^{n \times n}$ be an integral hyperbolic matrix, Ω its reduced automorphism group and $\mathcal{V}_i^{>0}$ and D as above. Then the residue class graph Γ_D/Ω is finite. In particular, the algorithm terminates.

Watson-Process

- ⇒ smaller determinant, less perfect points
- ⇒ acceleration of computation time

Watson-Process

- $\text{Aut}(L) \leq \text{Aut}(\text{Watson}(L)), [\text{Aut}(\text{Watson}(L)) : \text{Aut}(L)] < \infty$

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 - Compute $\text{Aut}(\text{Watson}(L))$.

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Watson-Process

- $\text{Aut}(L) \leq \text{Aut}(\text{Watson}(L))$, $[\text{Aut}(\text{Watson}(L)) : \text{Aut}(L)] < \infty$
- modified procedure:
 - Compute $\text{Aut}(\text{Watson}(L))$.
 - Find $\text{Aut}(L)$ via orbit-stabilizer calculation.

⇒ smaller determinant, less perfect points

⇒ acceleration of computation time

Statistics

	\emptyset Time (sec.)	\emptyset # Points
no Watson	70.08	7.4
Watson	2.96	2.7
Isotropic	98.18	8.9
Anisotropic	2.43	5.1

Table: Computational statistics

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$$A := \begin{pmatrix} 2 & 3 & 1 \\ 3 & 3 & 1 \\ 1 & 1 & 4 \end{pmatrix}, \quad A^{-1} = \frac{1}{11} \begin{pmatrix} -11 & 11 & 0 \\ 11 & -7 & -1 \\ 0 & -1 & 3 \end{pmatrix}$$

- D -perfect points:

$$\begin{aligned} x_1 &= (-1 \ 1 \ 0), & x_2 &= (-3 \ 3 \ 1), \\ x_3 &= (-4 \ 3 \ 2), & x_4 &= (-9 \ 11 \ 3). \end{aligned}$$

- stabilizers:

$$\text{Stab}(x_1) = \left\langle \begin{pmatrix} 1 & -2 & 0 \\ 0 & -1 & 0 \\ 0 & 0 & -1 \end{pmatrix} \right\rangle \cong C_2, \quad \text{Stab}(x_2) = \{1\}$$

$$\text{Stab}(x_3) = \left\langle \begin{pmatrix} 23 & -18 & -12 \\ -8 & 5 & 4 \\ 56 & -42 & -29 \end{pmatrix}, \begin{pmatrix} 23 & -18 & -12 \\ 0 & 1 & 0 \\ 44 & -36 & -23 \end{pmatrix} \right\rangle \cong C_2 \times C_2,$$

$$\text{Stab}(x_4) = \left\langle \begin{pmatrix} 5 & -8 & 0 \\ 3 & -5 & 0 \\ 1 & -2 & 1 \end{pmatrix} \right\rangle \cong C_2.$$

- stabilizers:

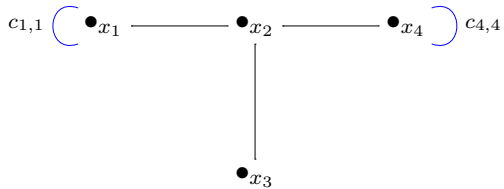
$$\text{Stab}(x_1) = \left\langle \begin{pmatrix} 1 & -2 & 0 \\ 0 & -1 & 0 \\ 0 & 0 & -1 \end{pmatrix} \right\rangle \cong C_2, \quad \text{Stab}(x_2) = \{1\}$$

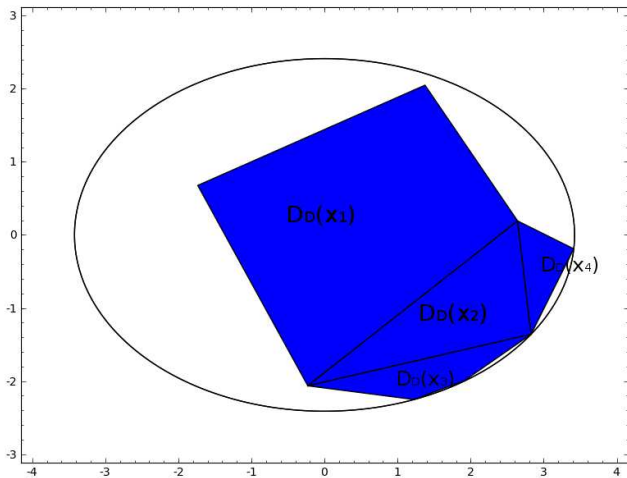
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$$\text{Stab}(x_4) = \left\langle \begin{pmatrix} 5 & -8 & 0 \\ 3 & -5 & 0 \\ 1 & -2 & 1 \end{pmatrix} \right\rangle \cong C_2.$$

- connecting elements:

$$c_{1,1} = \begin{pmatrix} 5 & -2 & 0 \\ 3 & -1 & 0 \\ 1 & 0 & -1 \end{pmatrix}, \quad c_{4,4} = \begin{pmatrix} 49 & -60 & -20 \\ 25 & -31 & -10 \\ 45 & -54 & -19 \end{pmatrix}.$$





An implementation of the algorithm in MAGMA as well as the slides of this talk are available on my homepage

<http://www.mi.uni-koeln.de/~mmertens>